

# CoLoSL

# Concurrent Local Subjective Logic

Azalea Raad

Jules Villard

Philippa Gardner

Imperial College London

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# Global Shared Resources

$P_1 \wedge P_2$

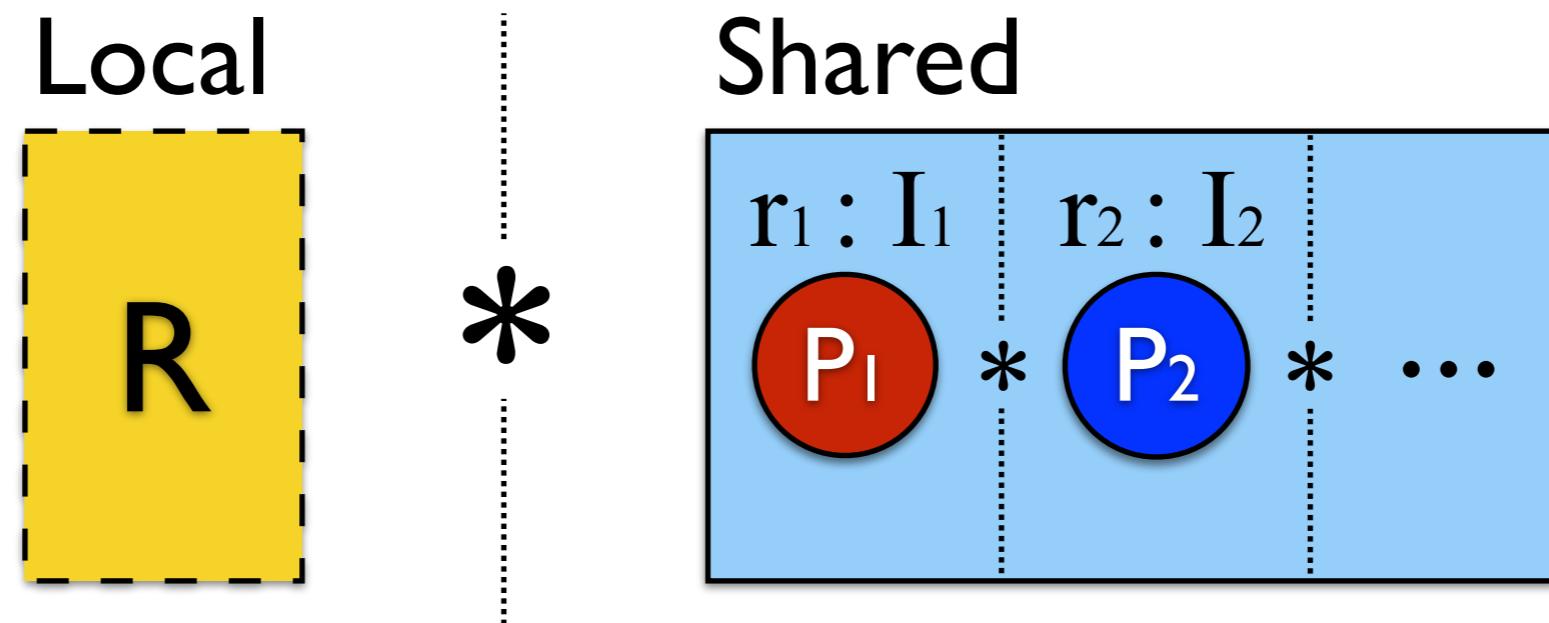
$$\frac{\left\{ \boxed{P_1} \right\} C_1 \left\{ \boxed{Q_1} \right\} \quad \left\{ \boxed{P_2} \right\} C_2 \left\{ \boxed{Q_2} \right\}}{\left\{ \boxed{P_1 \wedge P_2} \right\} C_1 \parallel C_2 \left\{ \boxed{Q_1 \wedge Q_2} \right\}}$$

# Global Shared Resources

$$\frac{\left\{ \boxed{P_1} \right\} C_1 \left\{ \boxed{Q_1} \right\} \quad \left\{ \boxed{P_2} \right\} C_2 \left\{ \boxed{Q_2} \right\}}{\left\{ \boxed{P_1 \wedge P_2} \right\} C_1 \parallel C_2 \left\{ \boxed{Q_1 \wedge Q_2} \right\}}$$

- ✿ No framing on shared resources / interference
  - ♦ Reasoning on GLOBAL resources
  - ♦ Interference on ALL resources considered
- ✿ No extension
  - ♦ cannot dynamically share resources/extend interference

# Disjoint Shared Resources



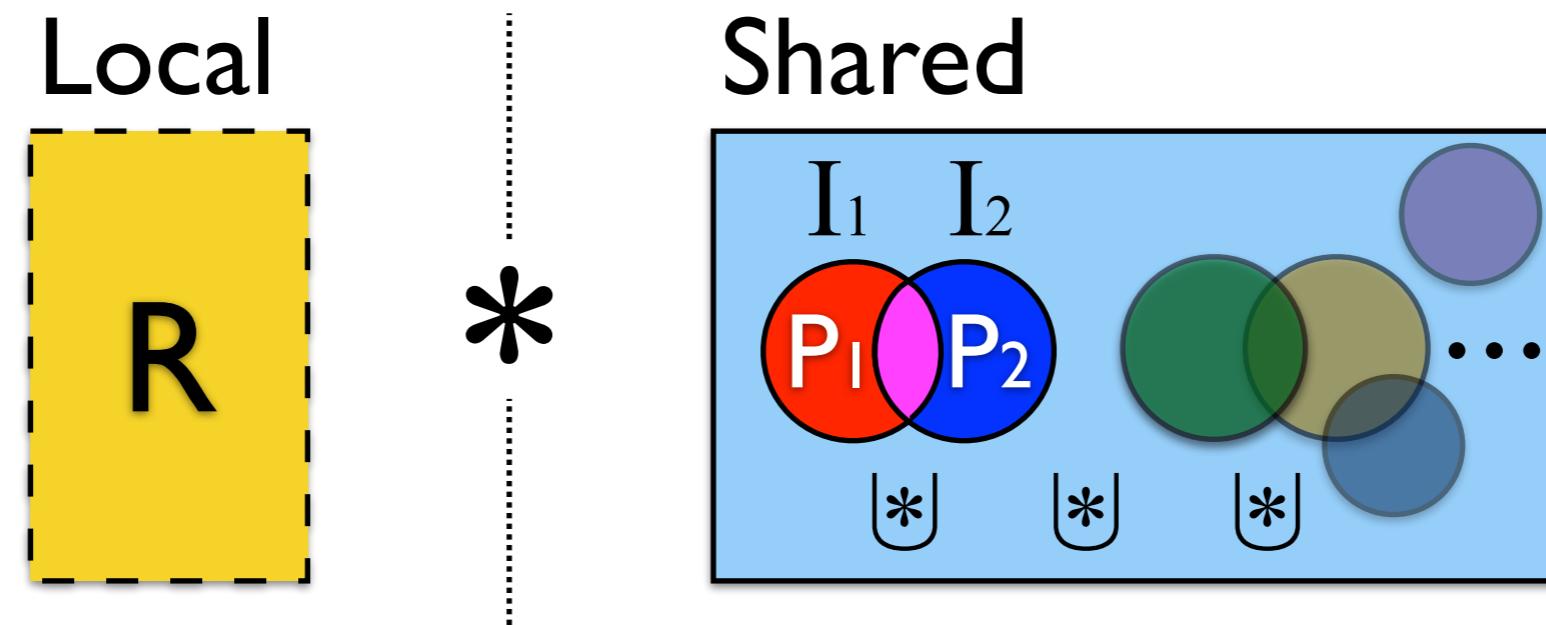
CSL, FCSL, CAP, HOCAP, iCAP, TaDA

# Disjoint Shared Resources

$$\frac{\left\{ \boxed{P_1}_{I_1} \right\} C_1 \left\{ \boxed{Q_1}_{I_1} \right\} \quad \left\{ \boxed{P_2}_{I_2} \right\} C_1 \left\{ \boxed{Q_2}_{I_2} \right\}}{\left\{ \boxed{P_1}_{I_1} * \boxed{P_2}_{I_2} \right\} C_1 \| C_2 \left\{ \boxed{Q_1}_{I_1} * \boxed{Q_2}_{I_2} \right\}}$$

- ✿ Limited framing on shared resources / interference
  - ♦ Static (pre-determined) frames (regions/ invariants)
  - ♦ Physically disjoint frames
- ✿ Limited extension
  - ♦ Can create new regions / invariants
  - ♦ Cannot extend regions with more resources/invariants

# CoLoSL: Concurrent Local Subjective Logic



CoLoSL

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$$\frac{\left\{ \boxed{P_1}_{I_1} \right\} C_1 \left\{ \boxed{Q_1}_{I_1} \right\} \quad \left\{ \boxed{P_2}_{I_2} \right\} C_1 \left\{ \boxed{Q_2}_{I_2} \right\}}{\left\{ \boxed{P_1}_{I_1} \uplus \boxed{P_2}_{I_2} \right\} C_1 \parallel C_2 \left\{ \boxed{Q_1}_{I_1} \uplus \boxed{Q_2}_{I_2} \right\}}$$

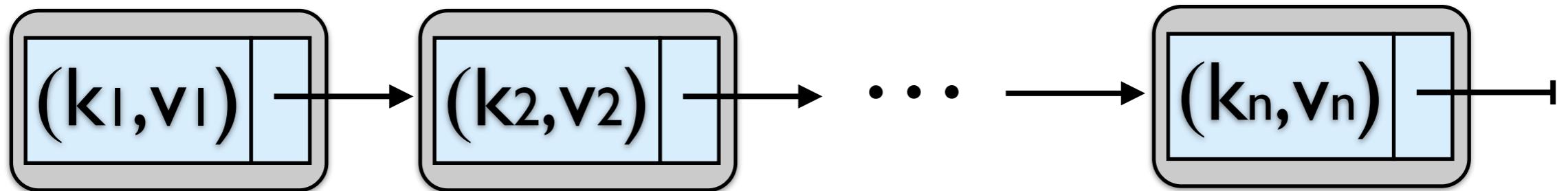
- ✿ Flexible framing on shared resources/invariants
  - ✧ Overlapping frames
  - ✧ Dynamic framing/rewriting of interference
- ✿ Dynamic extension
  - ✧ Extension of shared state with new resources/interference

# Examples

- ✿ B+ Tree (thanks to Shale Xiong)
  - ◆ Module composition; proof modularity; better abstraction
  - ◆ B+ Tree = List + BSTree

# Ordered Singly Linked-List

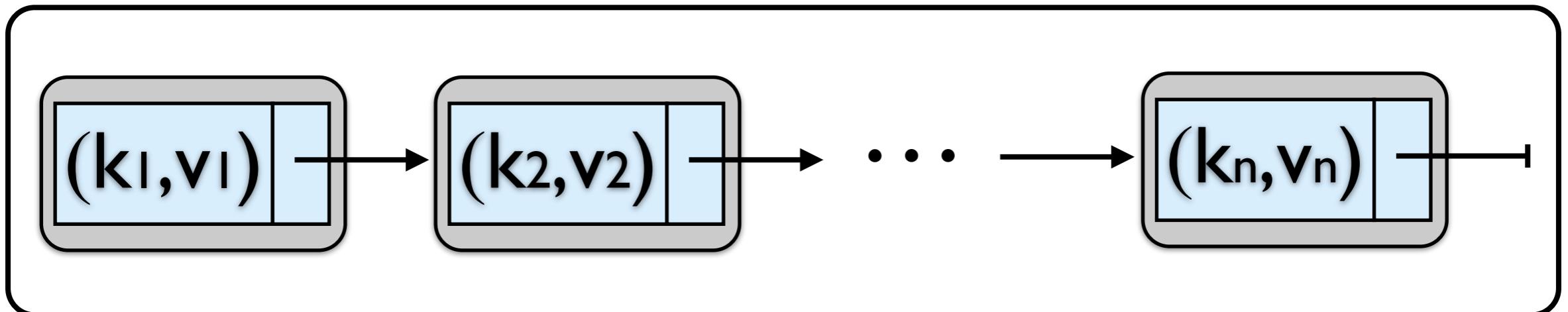
$\text{List}([(k_1, v_1), (k_2, v_2), \dots, (k_n, v_n)]) =$



- ✿ Value update
- ✿ Map (e.g. increment all elements)
- ✿ Insertion (involves pointer surgery)
- ✿ Removal (involves pointer surgery)

# Concurrent Ordered Singly Linked-List

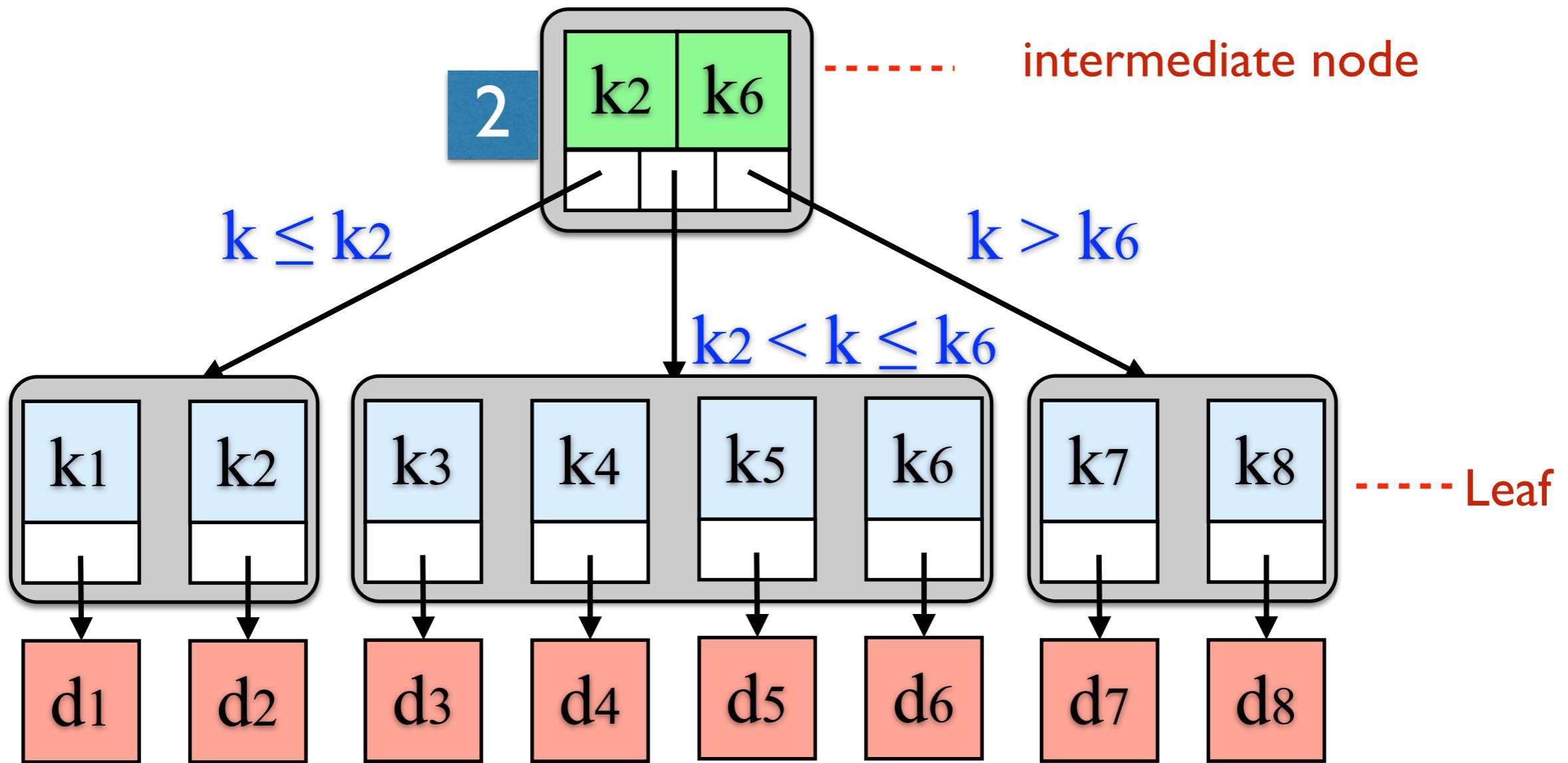
$\text{List}([(k_1, v_1), (k_2, v_2), \dots, (k_n, v_n)]) =$



$$I_L = I_{\text{up}} \cup I_{\text{map}} \cup I_{\text{add}} \cup I_{\text{rem}}$$

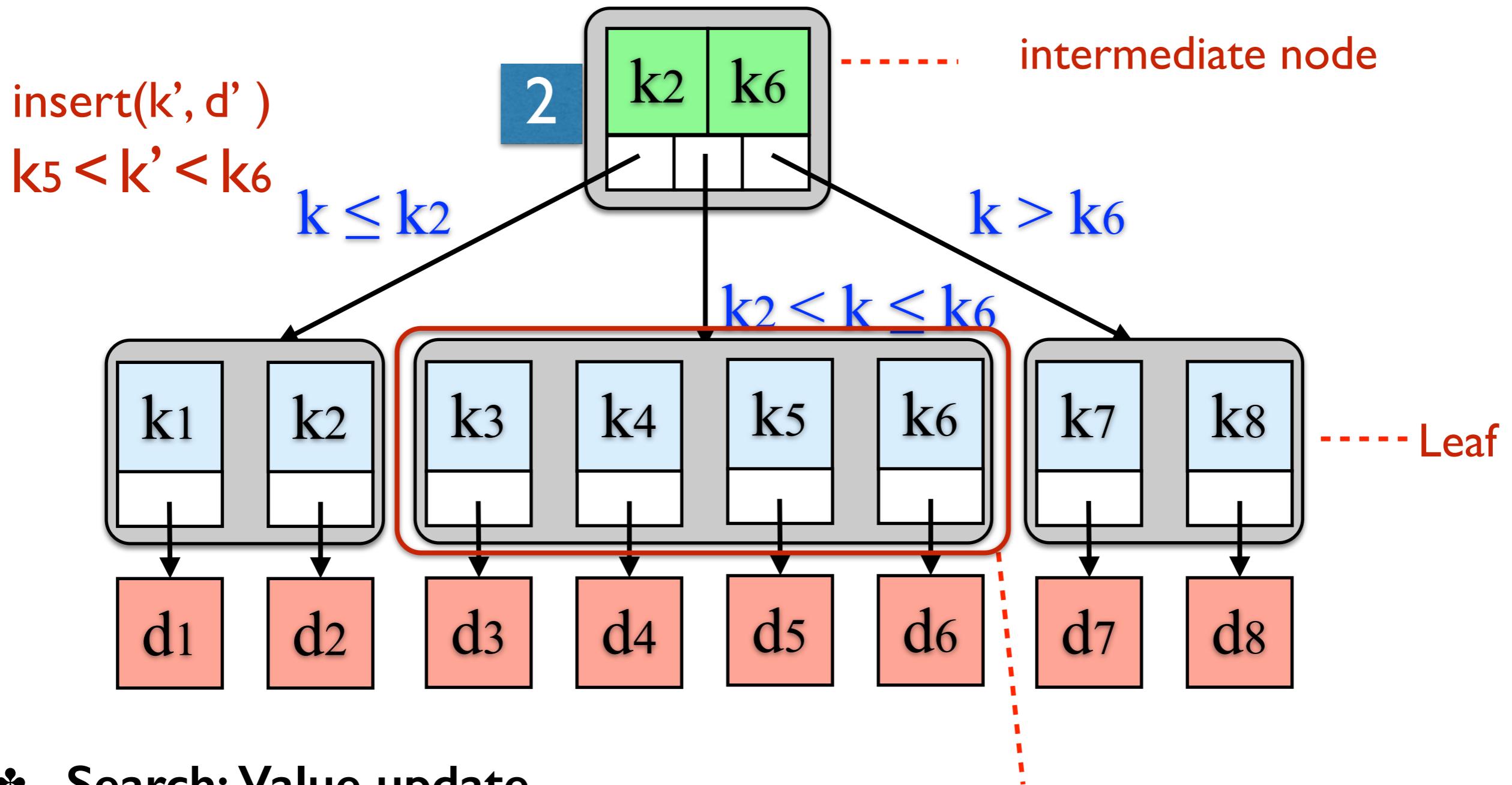
- ✿ Value update
- ✿ Map (e.g. increment all elements)
- ✿ Insertion (involves pointer surgery)
- ✿ Removal (involves pointer surgery)

# Balanced Search Tree (Degree 2)



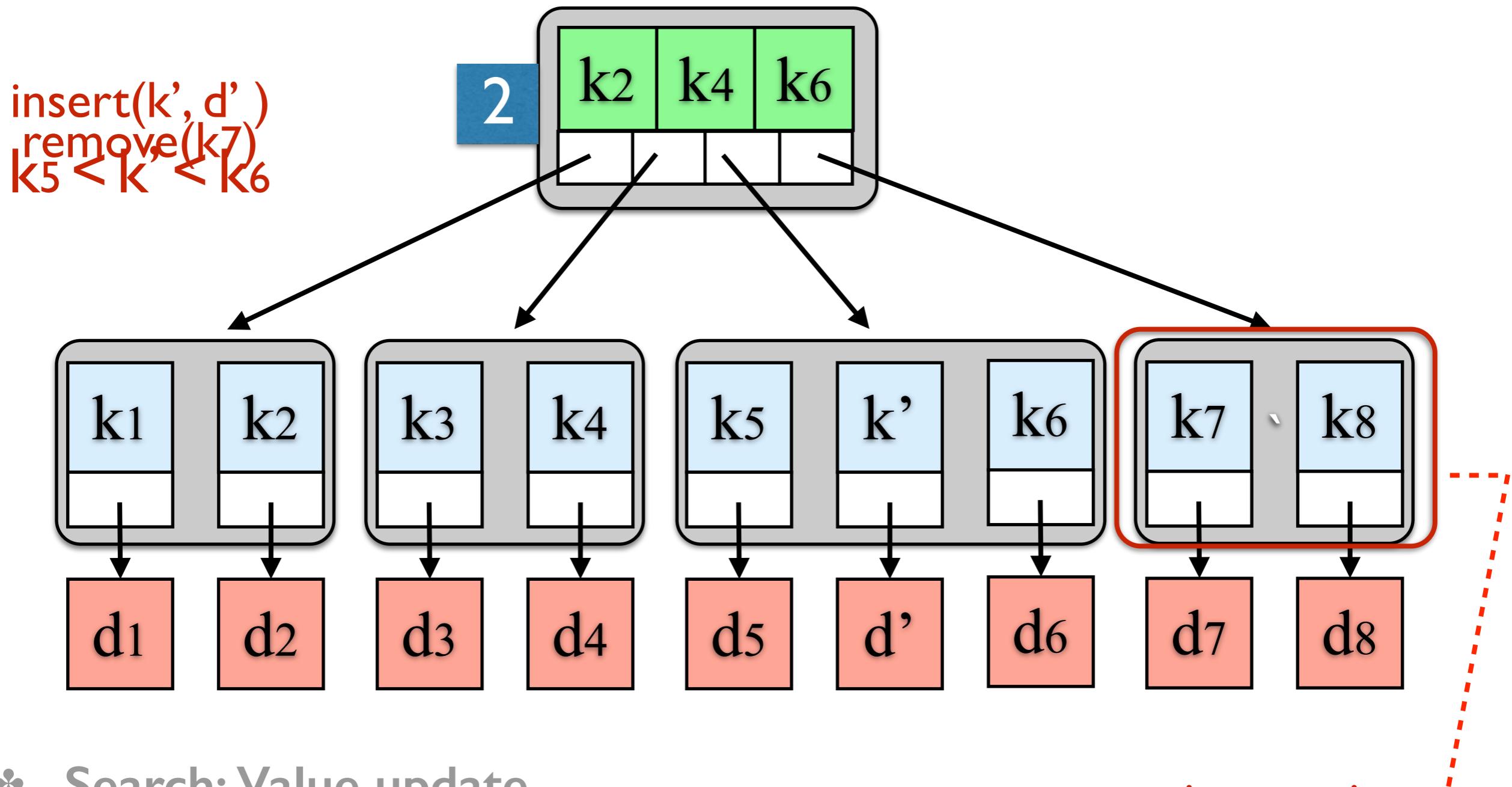
- Balanced: all immediate subtrees of a node have the same height
- Leaf-heavy: data (values) stored in leaf nodes
- Degree (d): no. of children (m) on each node  $d \leq m \leq 2d$

# Balanced Search Tree (Degree 2)



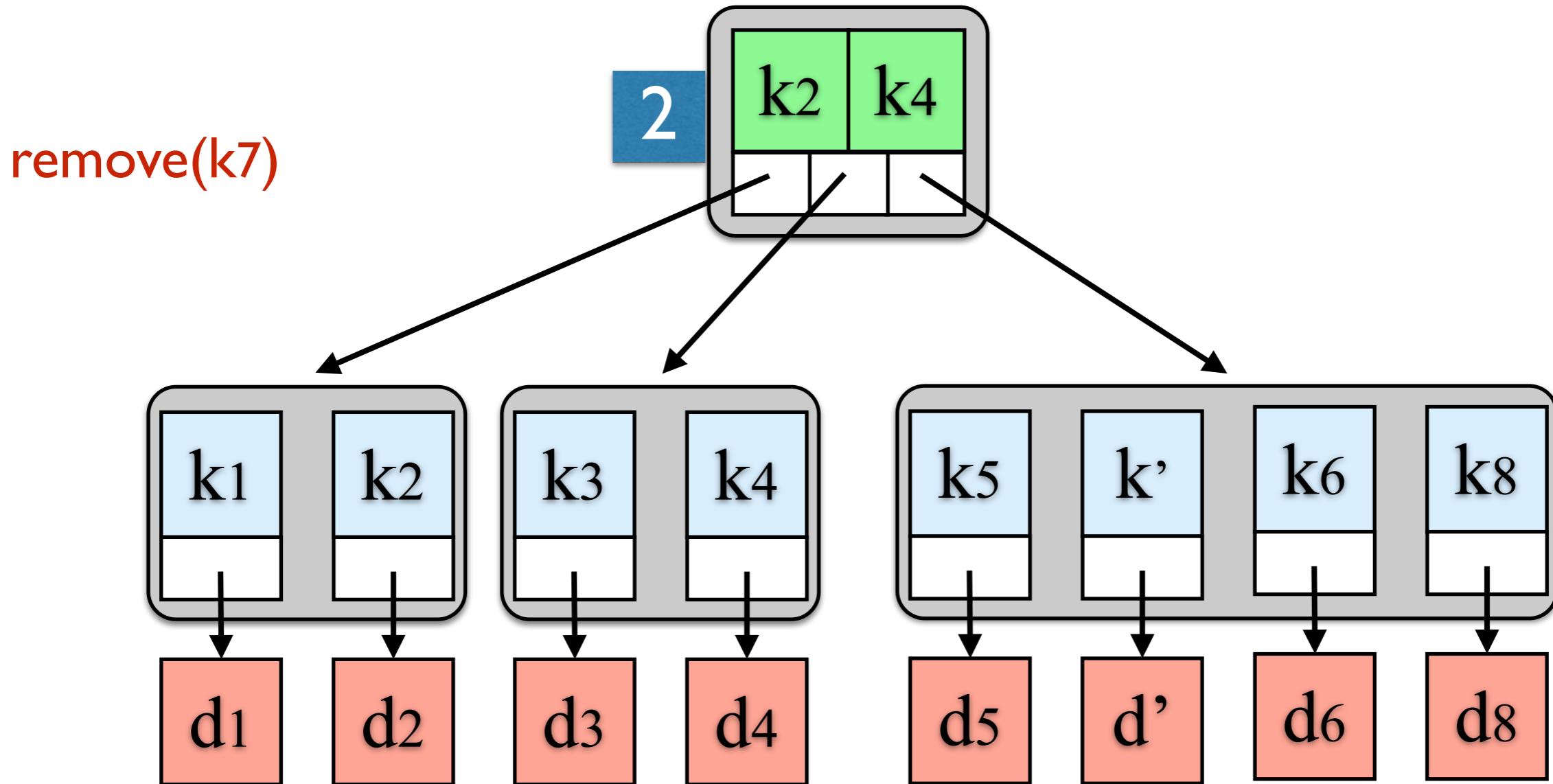
- ⌘ Search; Value update
- ⌘ Insertion may require splitting at max capacity
- ⌘ Removal

# Balanced Search Tree



- ⌘ Search; Value update at min capacity
- ⌘ Insertion may require splitting
- ⌘ Removal may require merging

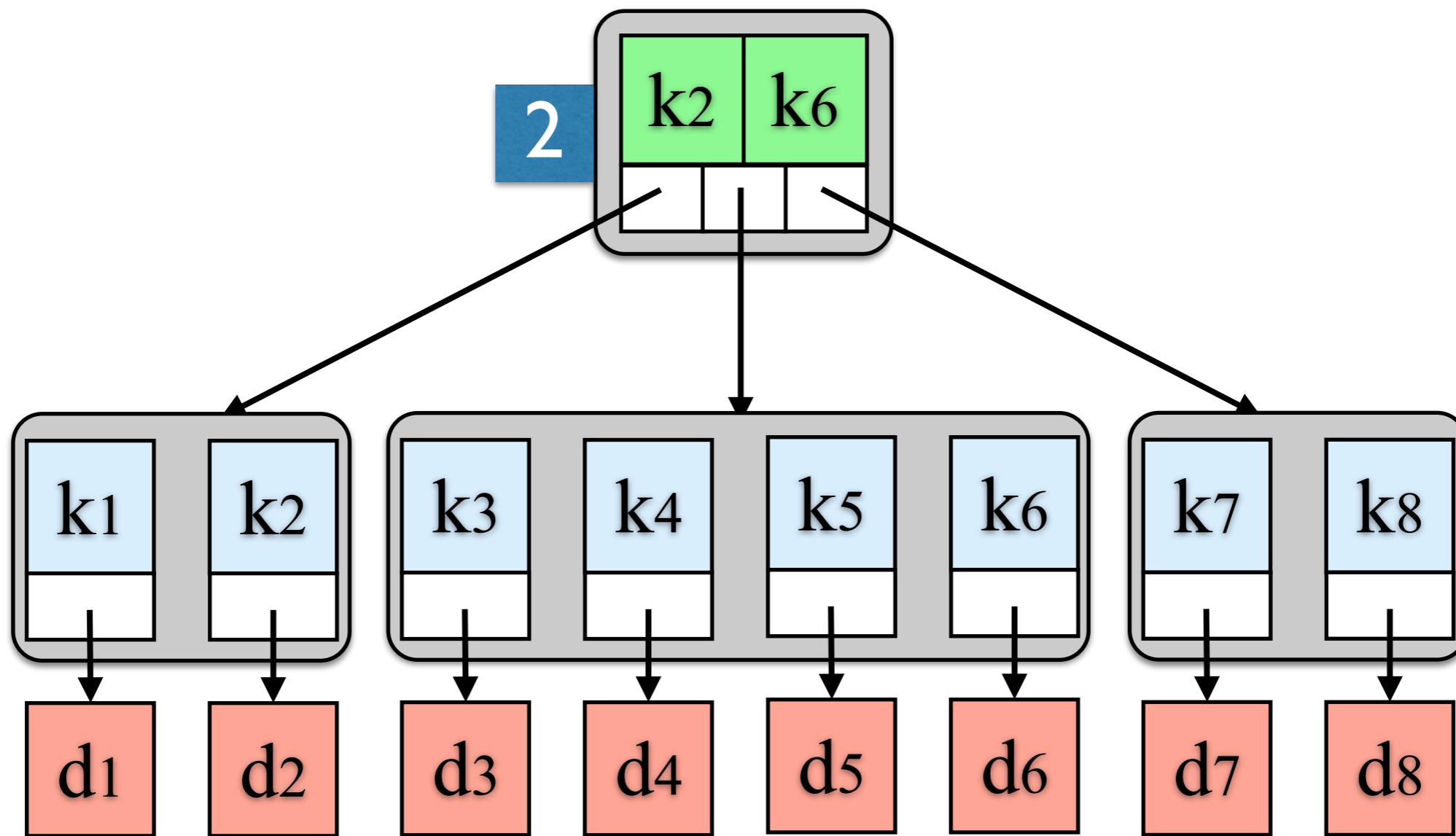
# Balanced Search Tree



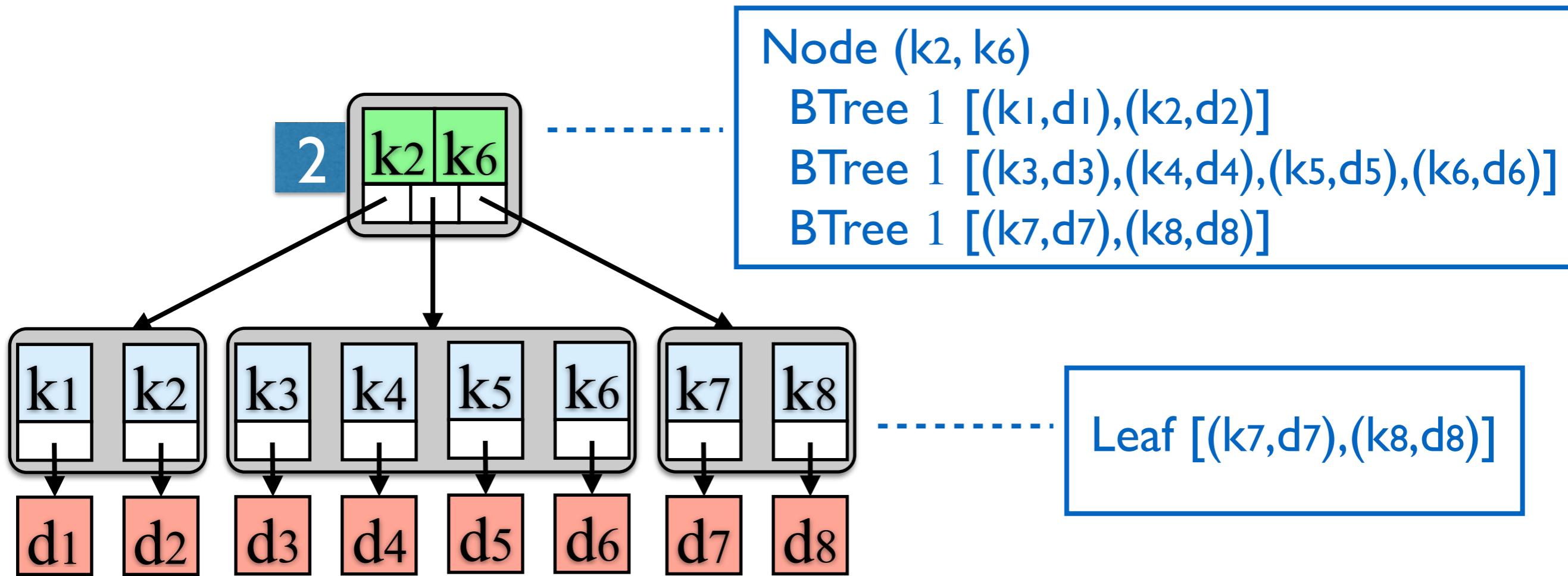
- ⌘ Search; Value update
- ⌘ Insertion
- ⌘ Removal may require merging

# Balanced Search Tree

BSTree 2  $\left( \left[ \left( k_1, d_1 \right) \dots \left( k_8, d_8 \right) \right] \right) =$



# Balanced Search Tree



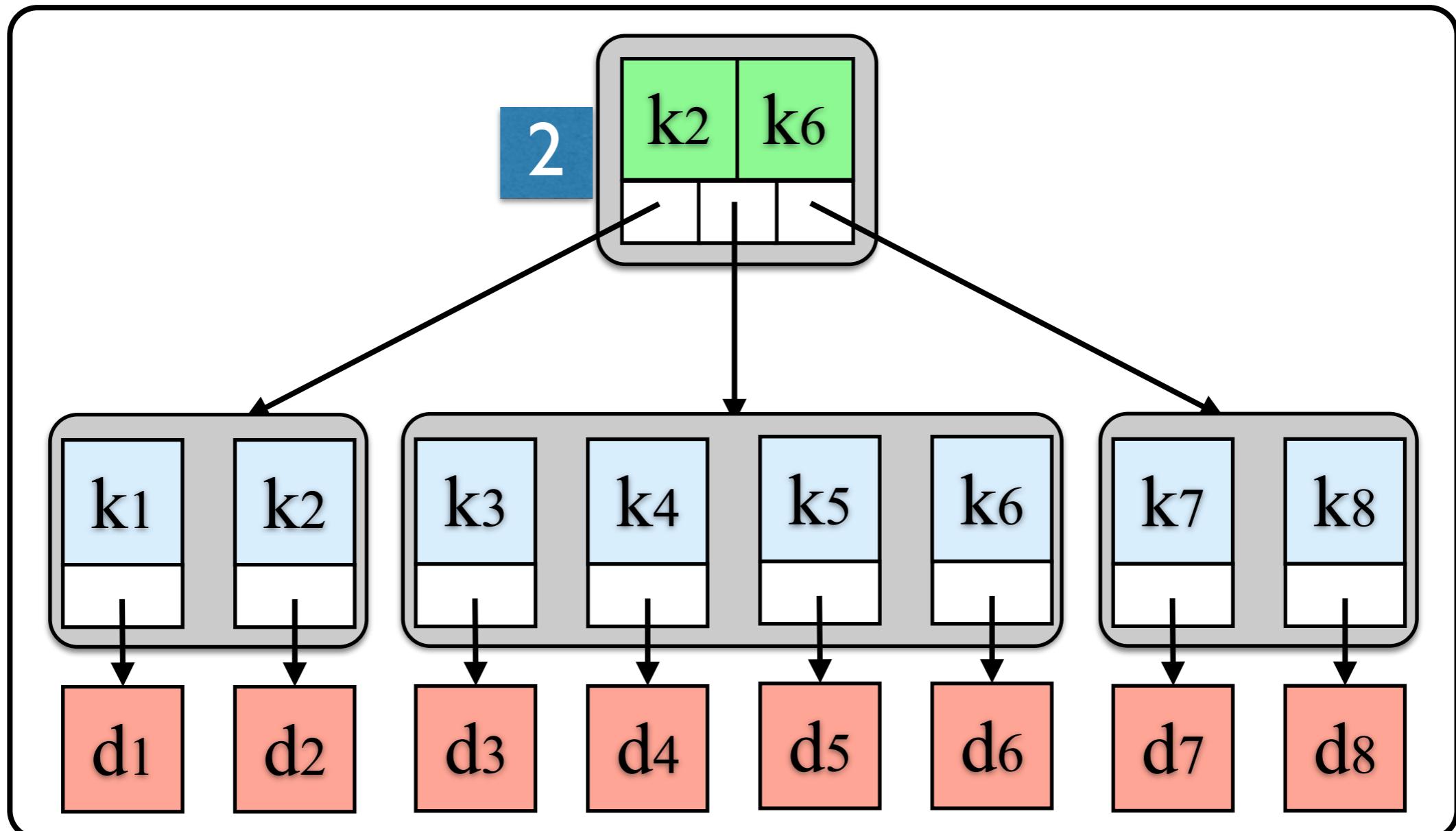
$\text{BSTree } (h+1) \ L =$

$\exists k_1, \dots, k_m, L_1, \dots, L_{(m+1)} . \ L = L_1 \cup \dots \cup L_{(m+1)}$   
Node  $(k_1, \dots, k_m)$  (BSTree  $h L_1$ ) ... (BSTree  $h L_{(m+1)}$ )

$\text{BSTree } 1 \ L = \text{Leaf } L$

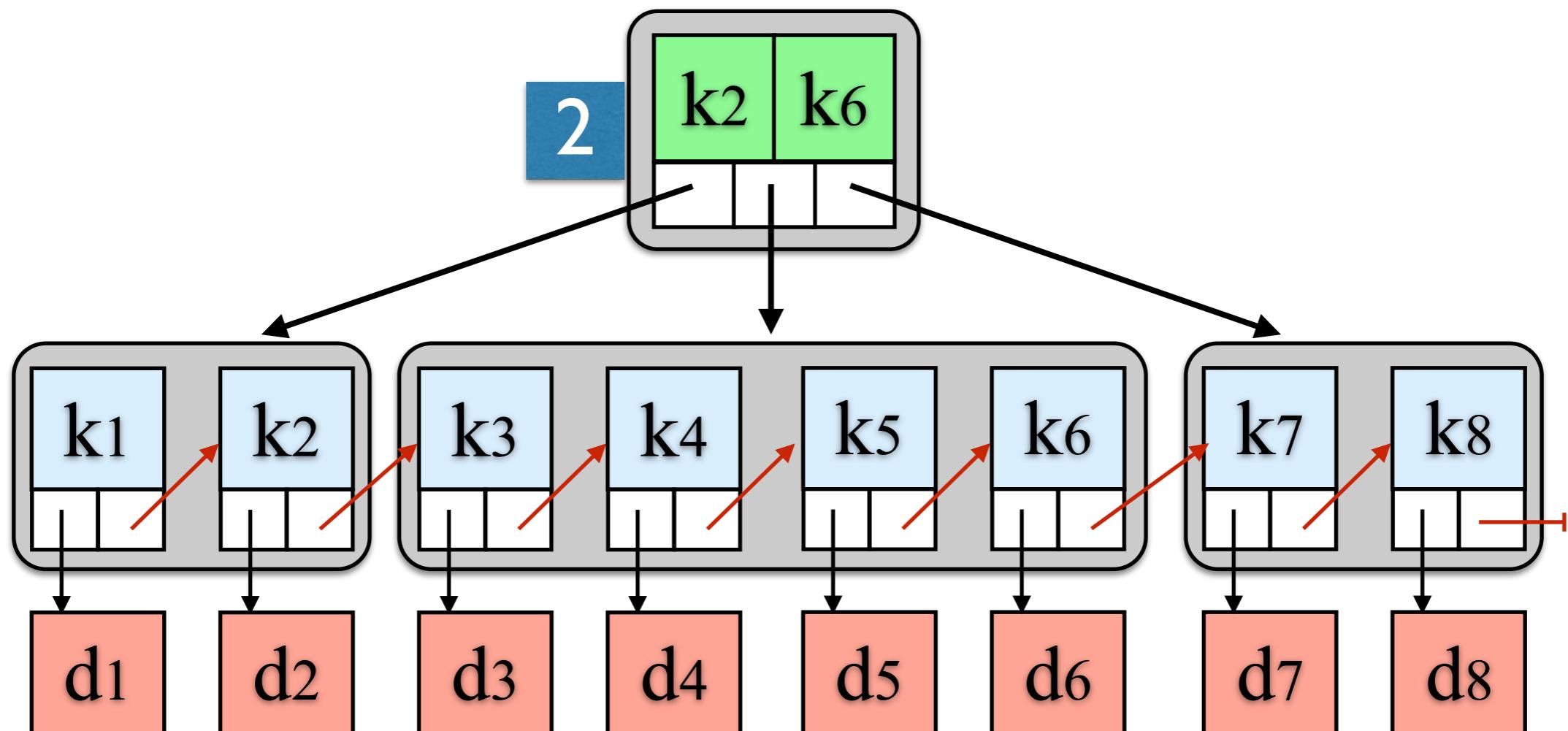
# Concurrent Balanced Search Tree

$\text{BSTree}([(k_1, d_1) \dots (k_8, d_8)]) =$



$$I_B = I_{\text{up}} \cup I_{\text{find}} \cup I_{\text{add}} \cup I_{\text{rem}}$$

# B+ Tree



# B+ Tree

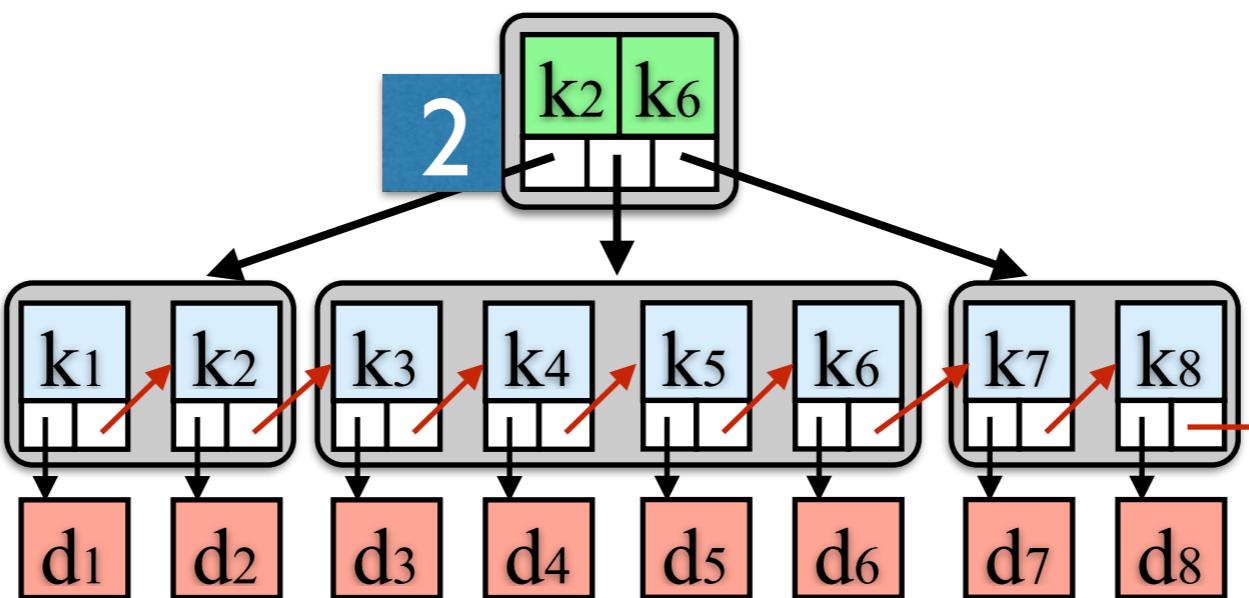
B+ Tree h L

$\Leftrightarrow$

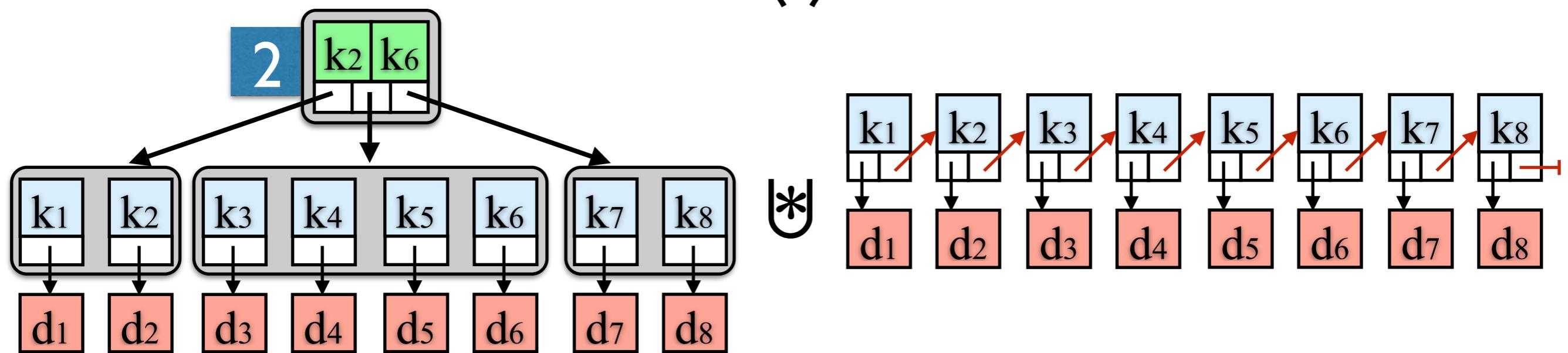
BSTree h L

$\bowtie$

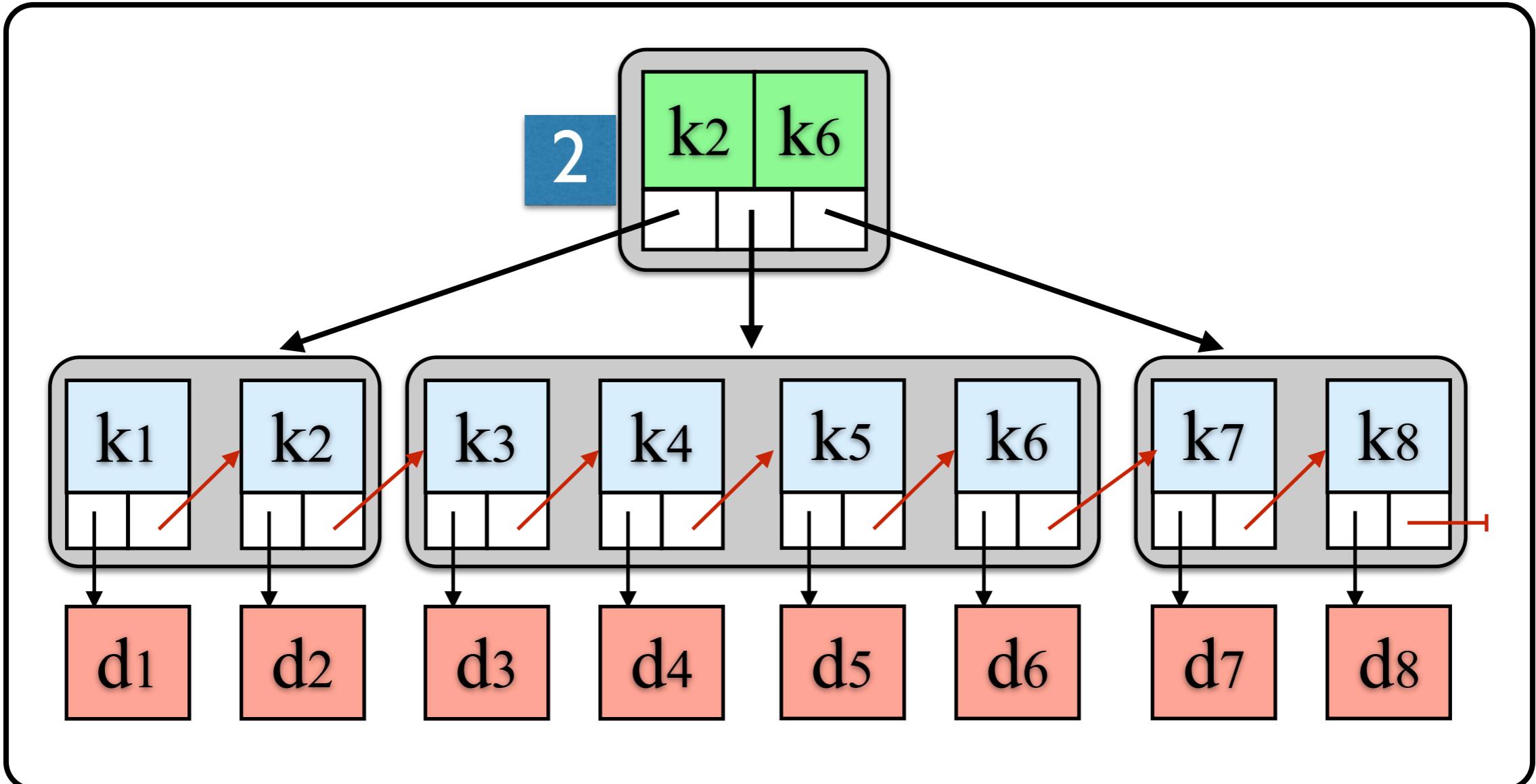
List L



$\Leftrightarrow$



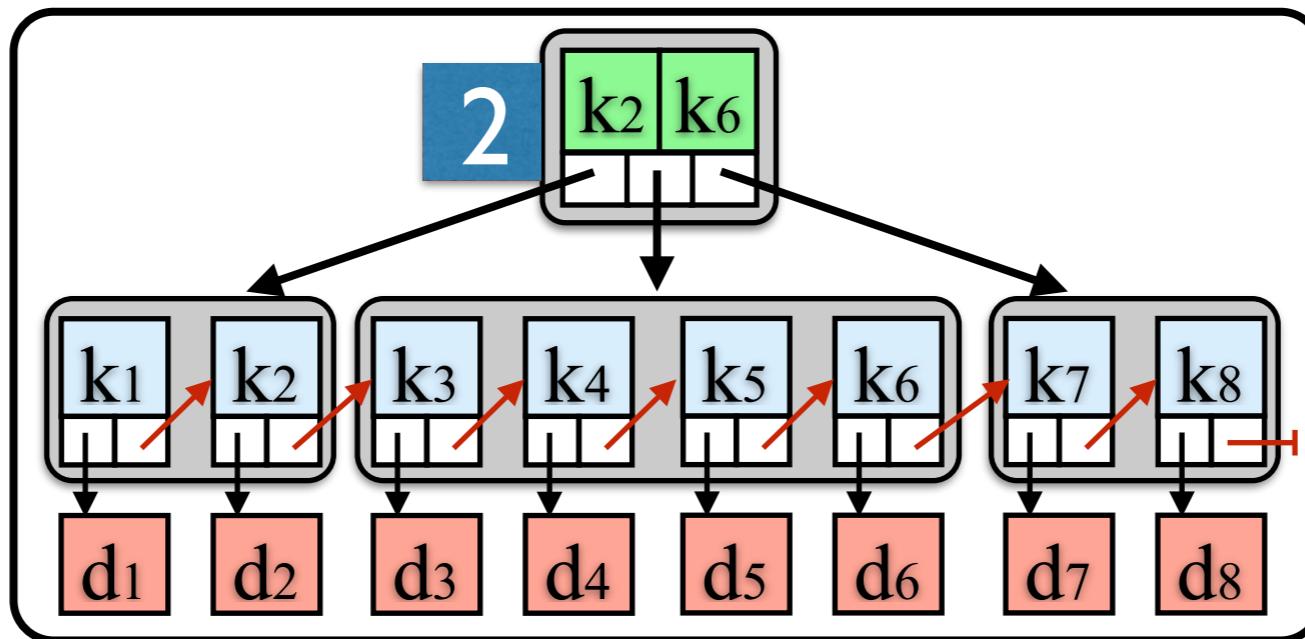
# Concurrent B+ Tree



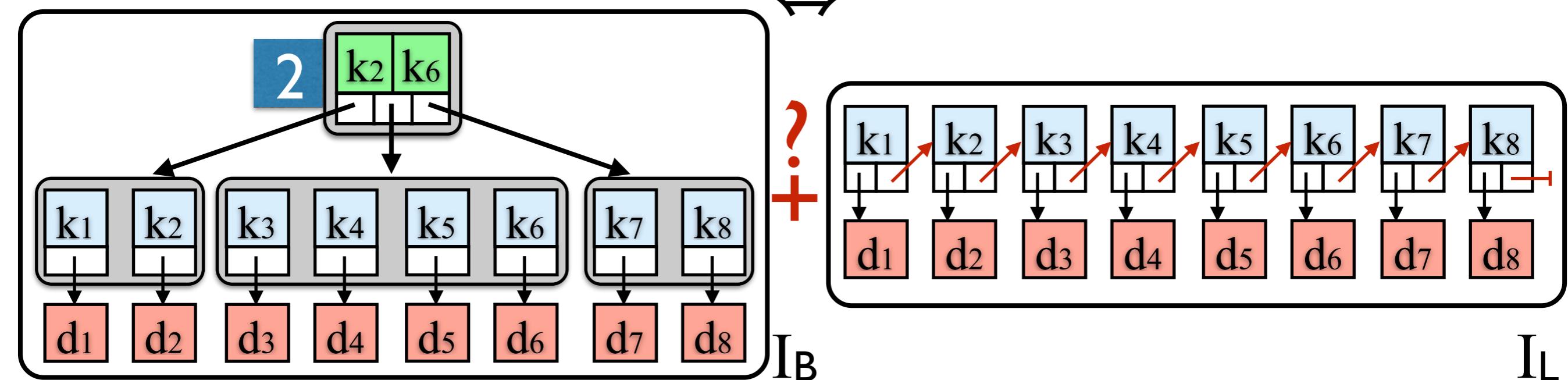
$$I = I_B \cup I_L$$

# Concurrent B+ Tree Wish List

$$\boxed{\text{B+ Tree } h \ L}_{I_{BUI_L}} \Leftrightarrow \boxed{\text{BSTree } h \ L}_{I_B} + \boxed{\text{List } L}_{I_L}$$

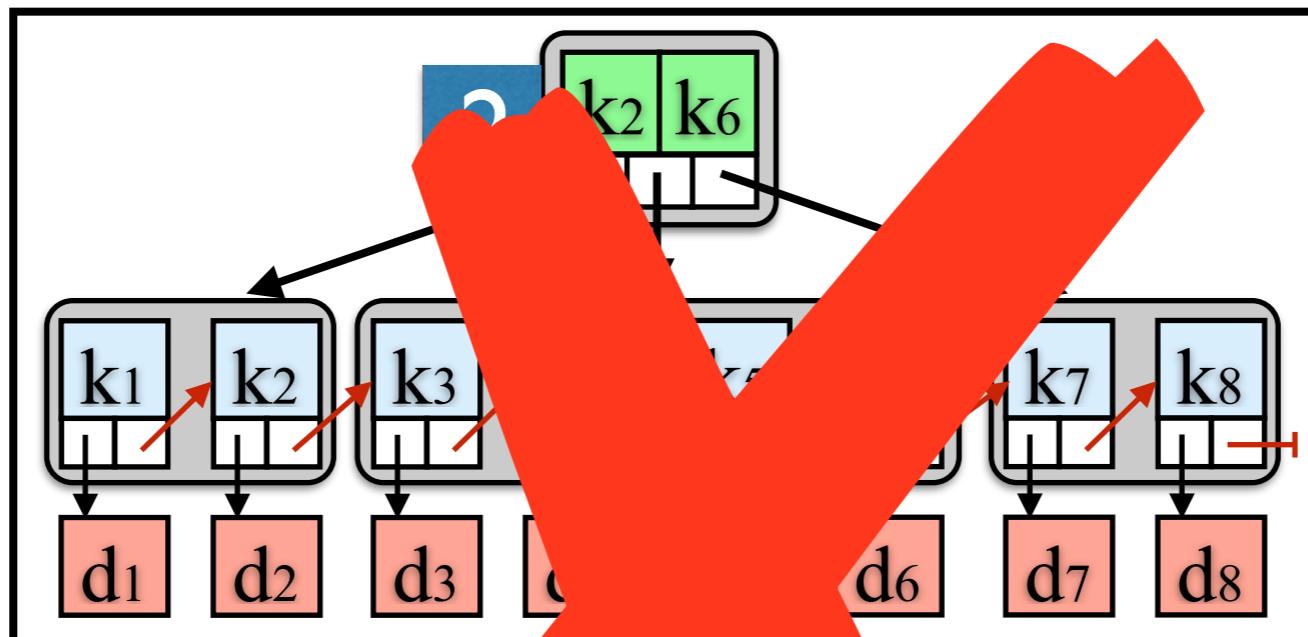


$\Leftrightarrow$

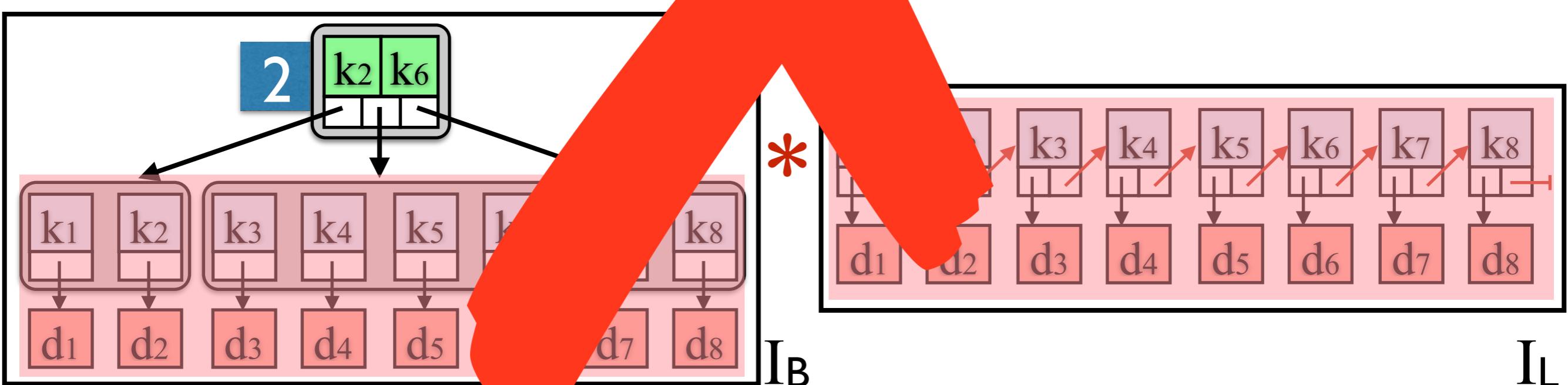


# Concurrent B+ Tree (Existing Approaches)

$$\boxed{\text{B+ Tree } h \ L}_{I_{\text{BUIL}}} \Leftrightarrow \boxed{\text{BSTree } h \ L}_{I_{\text{B}}} * \boxed{\text{List } L}_{I_{\text{L}}}$$

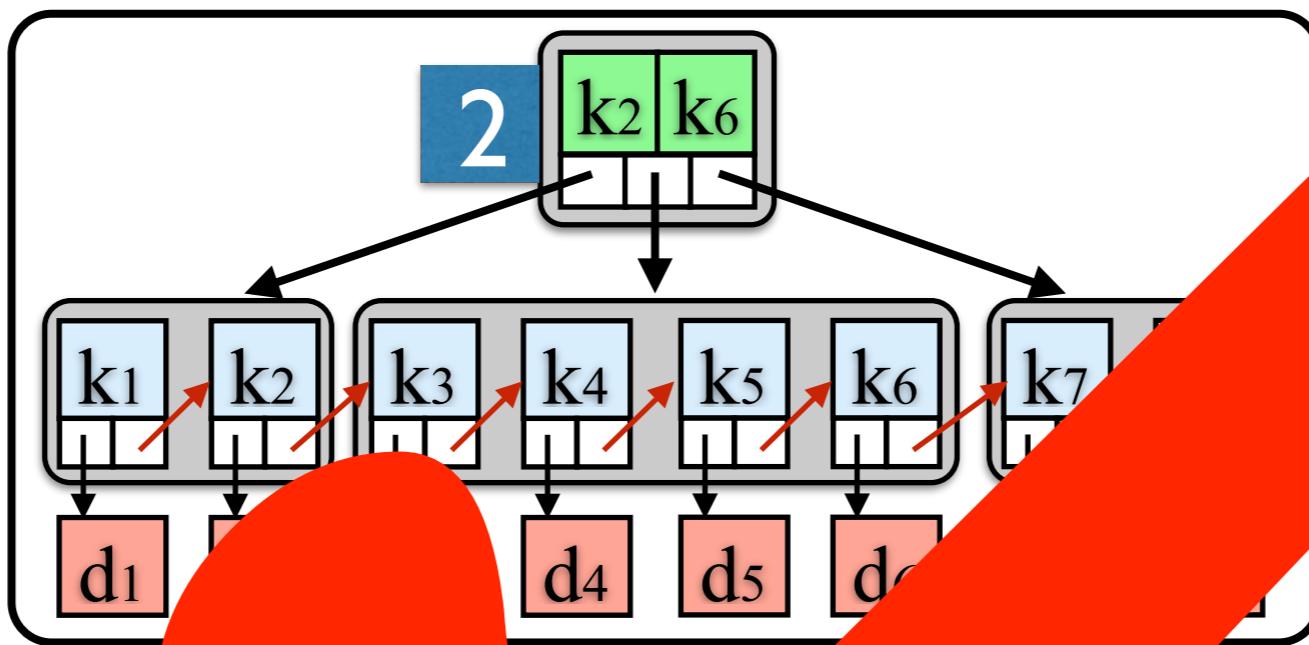


$$I = I_B \cup I_L$$



# Concurrent B+ Tree (CoLoSL)

$$\boxed{\text{B+ Tree } h \ L}_{I_{BUI_L}} \Leftrightarrow \boxed{\text{BSTree } h \ L}_{I_B} \textcolor{red}{\heartsuit} \boxed{\text{List } L}_{I_L}$$



$$I = I_B \cup I_L$$



# B+ Tree Wish List

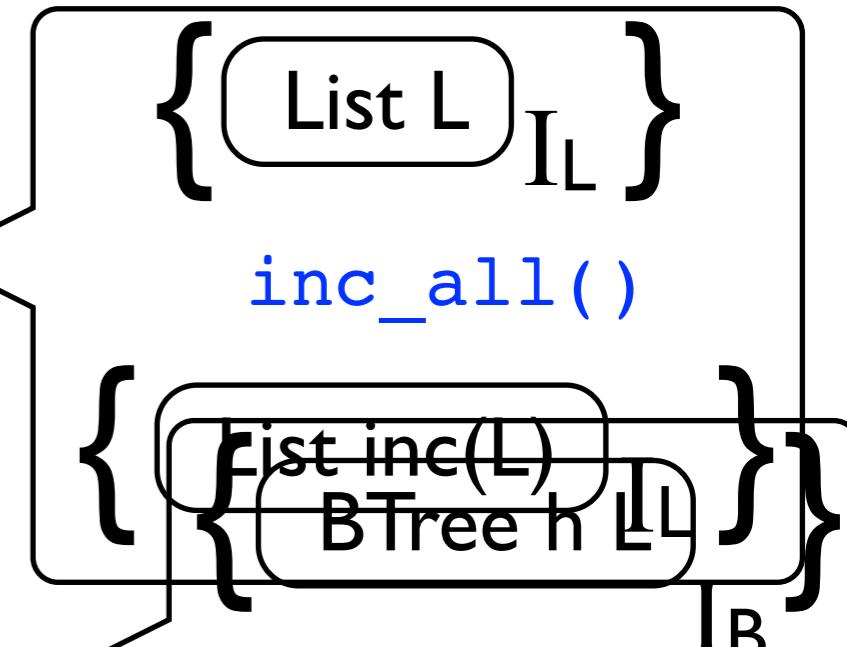
- ❖ Module Composition

$$\boxed{\text{B+ Tree h L}}_{I_{BUIL}} \Leftrightarrow \boxed{\text{BTree h L}}_{I_B} \between \boxed{\text{List L}}_{I_L}$$

- ❖ Proof Modularity (overlapping frames)

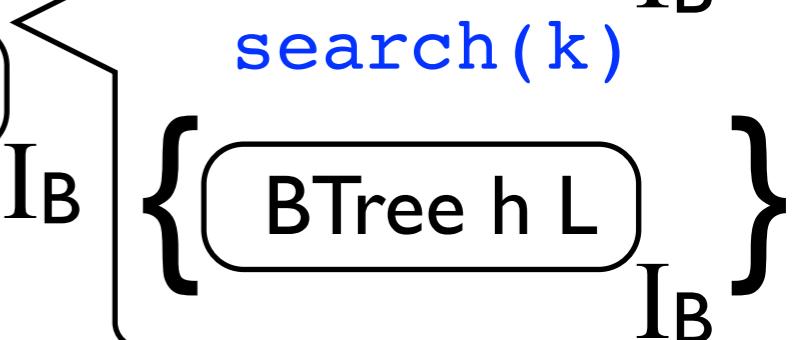
- ❖ Reuse list-only operation (e.g. map)

$$\boxed{\text{BTree h L}}_{I_B} \between \boxed{\text{List L}}_{I_L} \xrightarrow{\text{(frame)}} \boxed{\text{List L}}_{I_L}$$



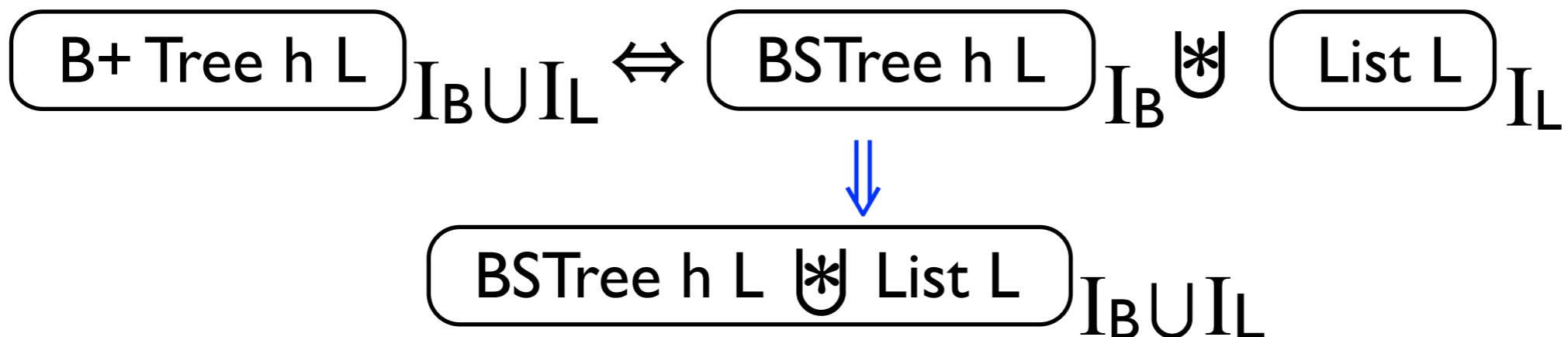
- ❖ Reuse tree-only operation (e.g. search)

$$\boxed{\text{BTree h L}}_{I_B} \between \boxed{\text{List L}}_{I_L} \xrightarrow{\text{(frame)}} \boxed{\text{BTree h L}}_{I_B}$$



- ❖ Combine tree/list module operations to implement B+ tree operations (e.g. remove)

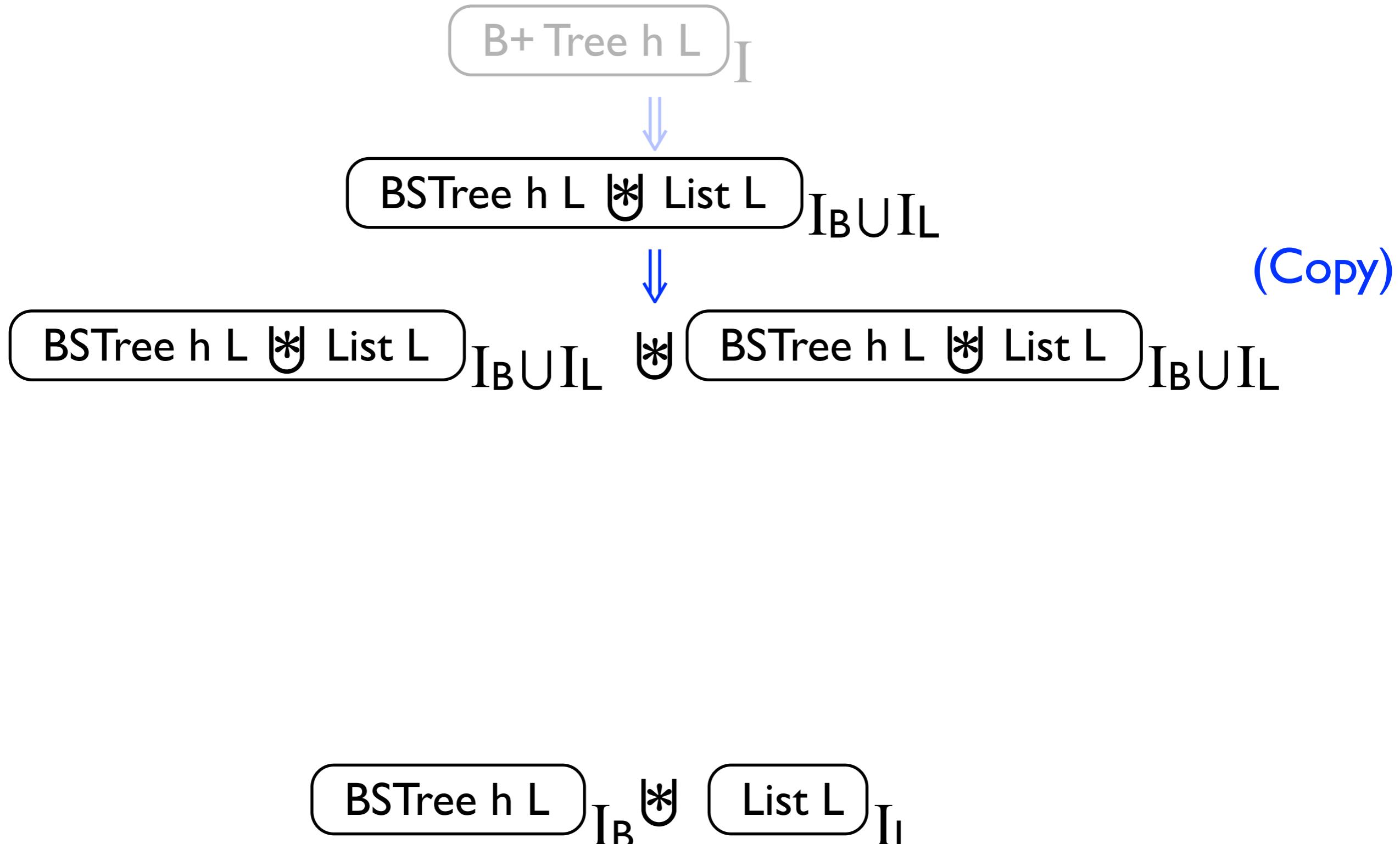
# CoLoSL Principles



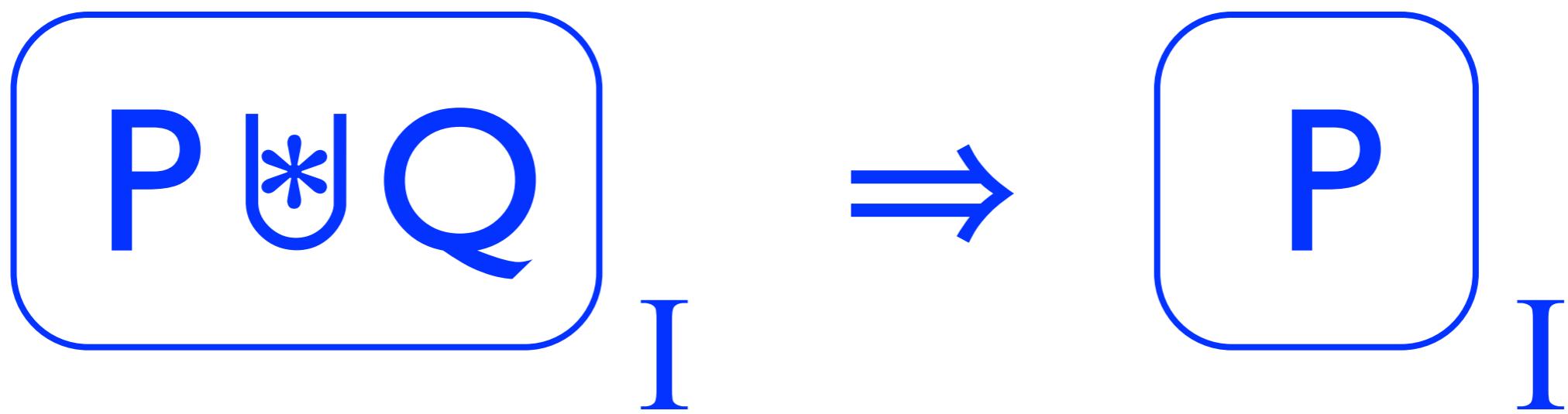
# Duplicating Resources



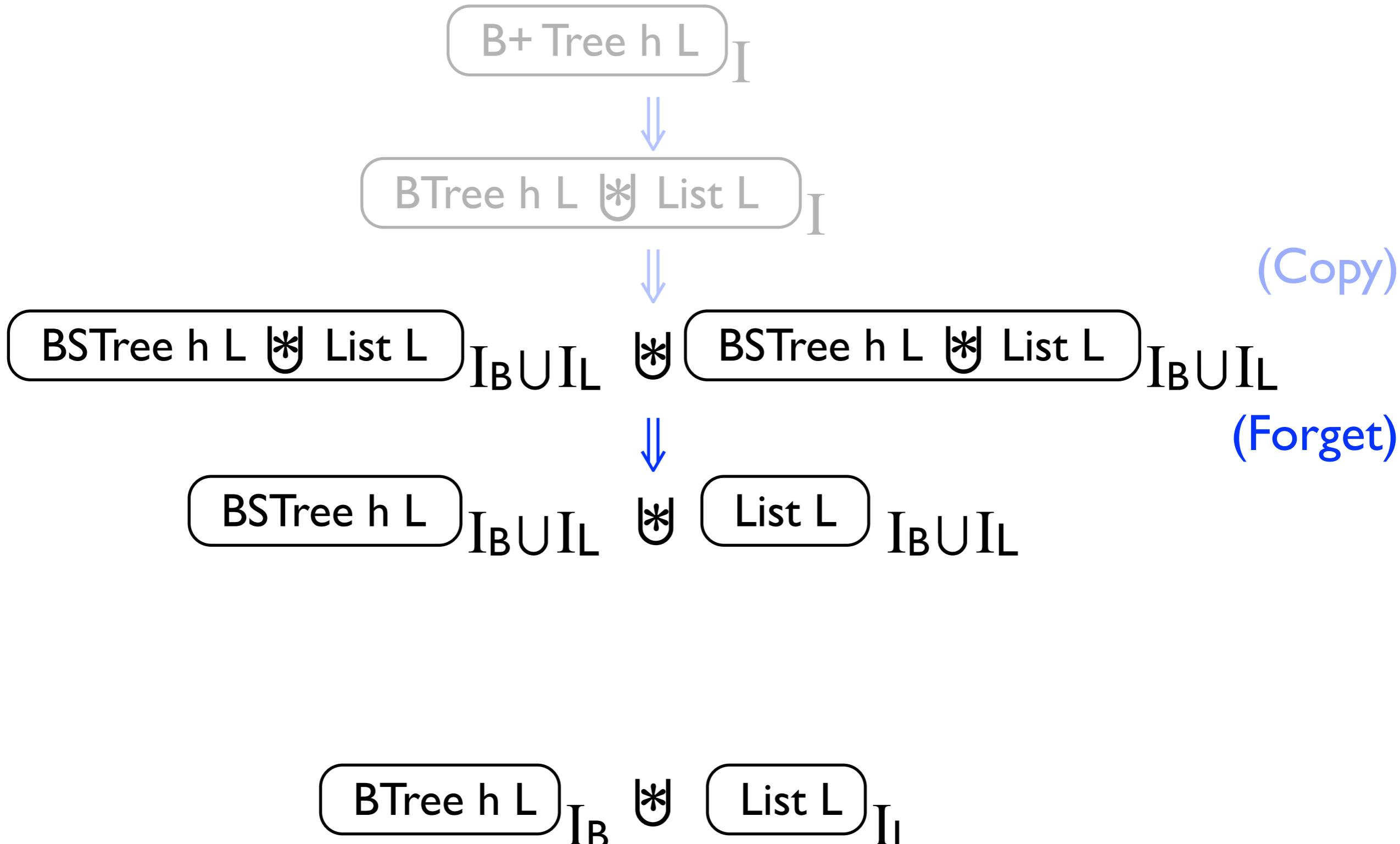
# CoLoSL Principles



# Forgetting Resources



# CoLoSL Principles

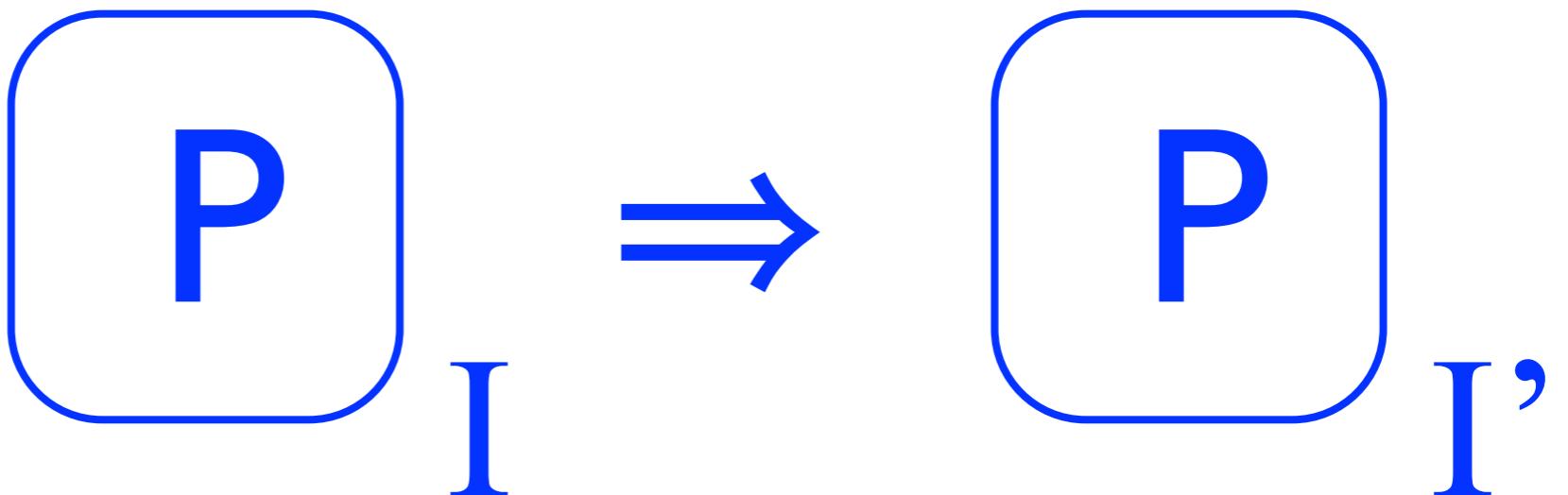


# Forgetting Interference (Shift)

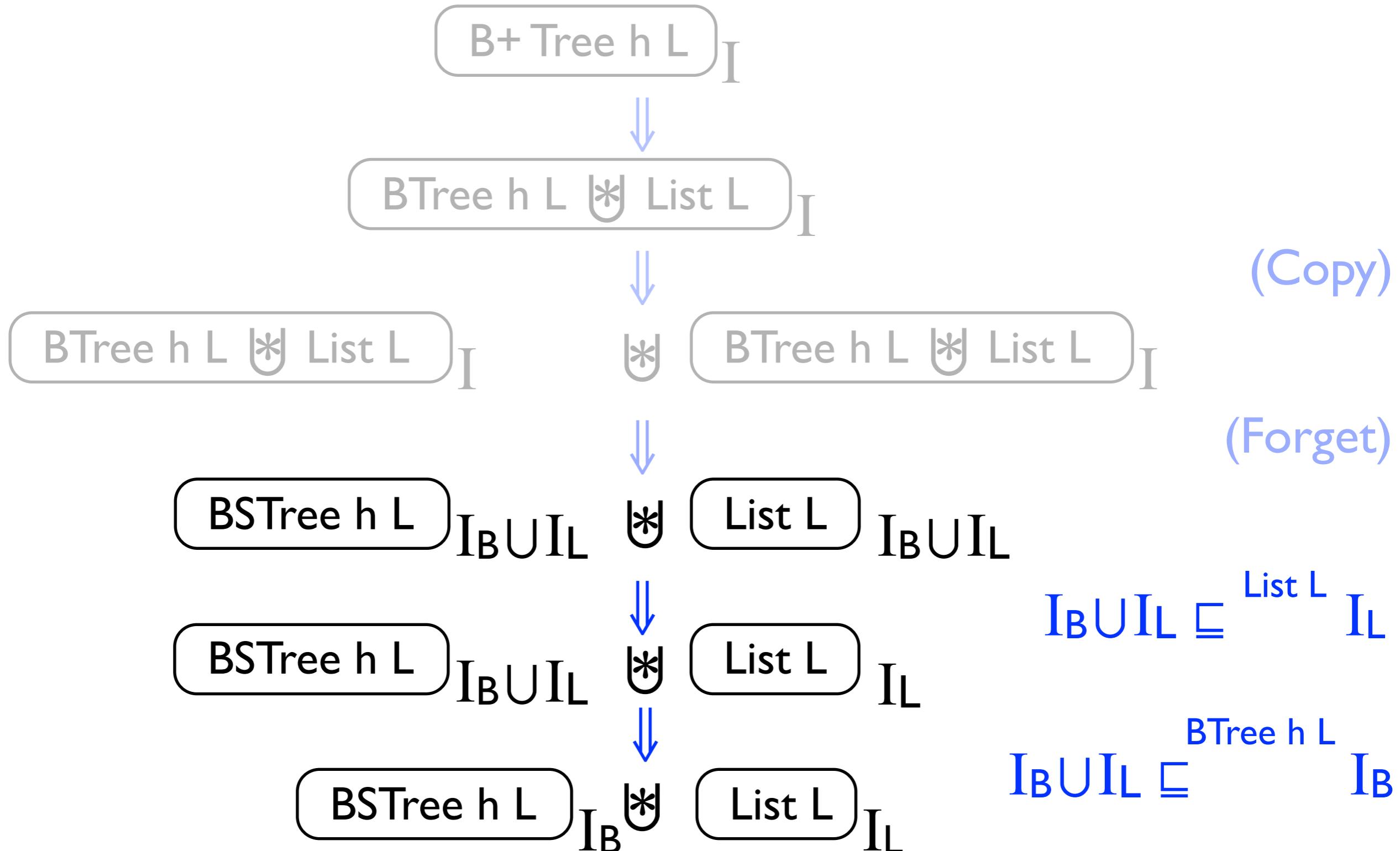
if

$$I \sqsubseteq^P I'$$

then



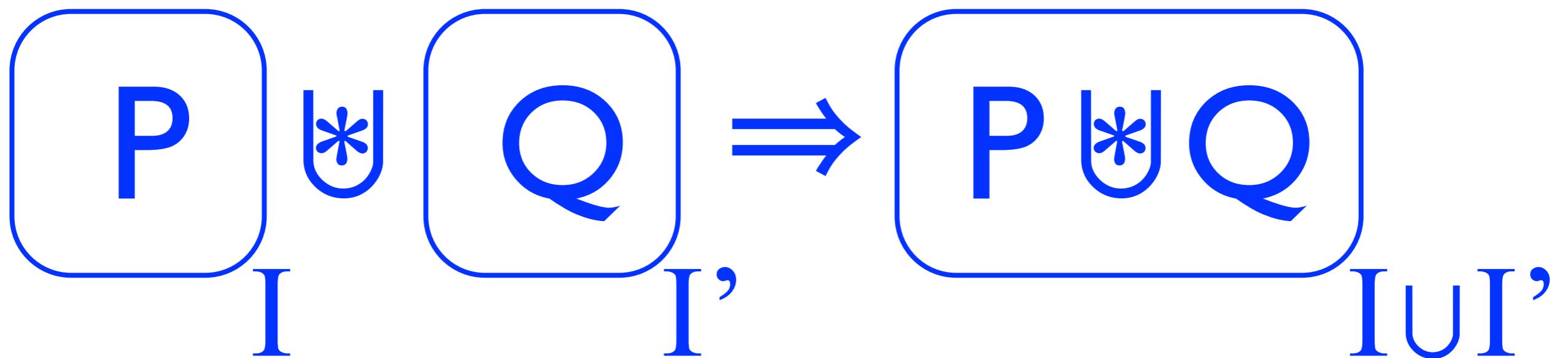
# CoLoSL Principles



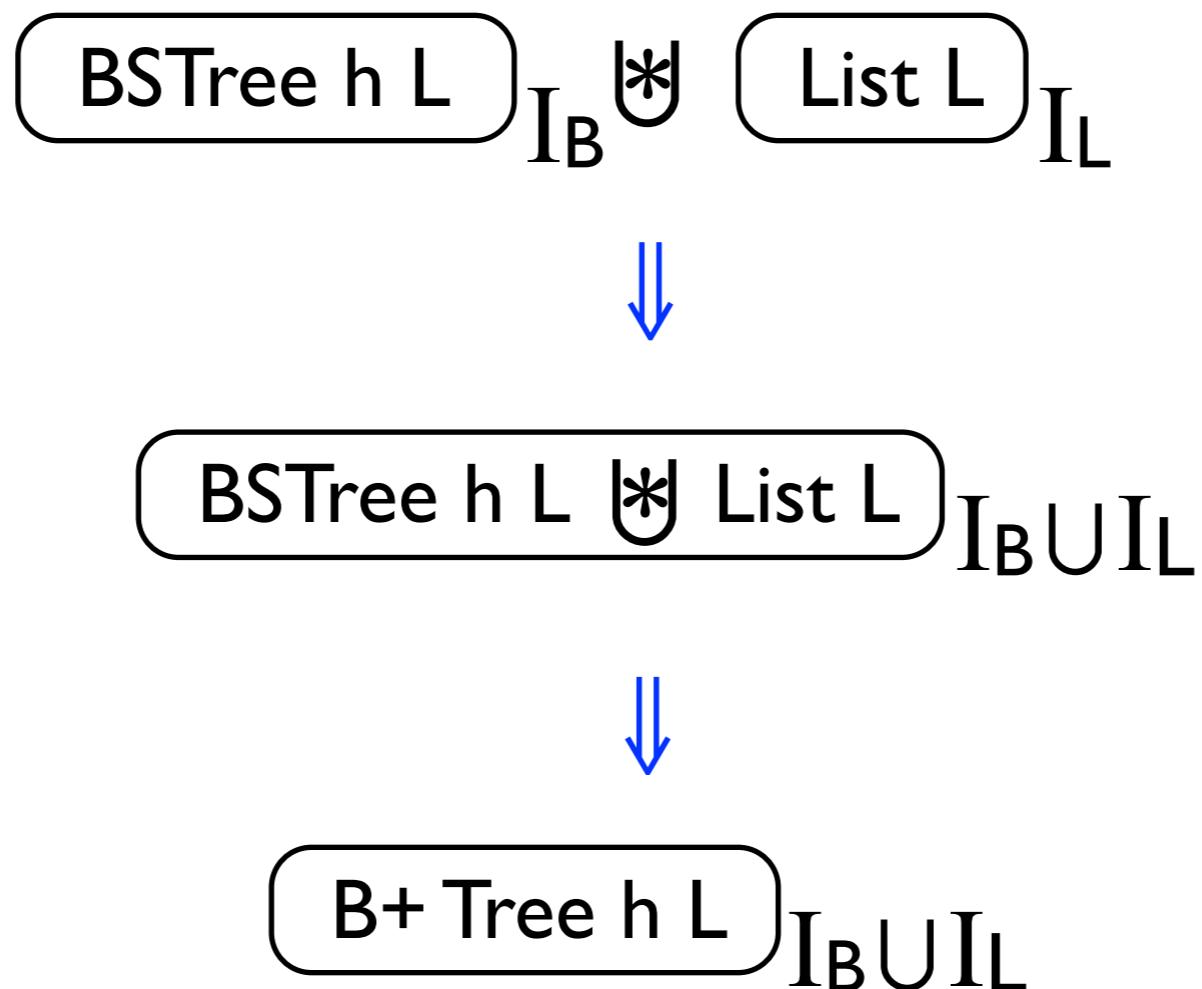
# CoLoSL Principles

$$\boxed{\text{B+ Tree } h \ L}_{I_{BUIL}} \Rightarrow \boxed{\text{BSTree } h \ L}_{I_B} \bowtie \boxed{\text{List } L}_{I_L}$$
$$\boxed{\text{B+ Tree } h \ L}_{I_{BUIL}} \Leftarrow \boxed{\text{BSTree } h \ L}_{I_B} \bowtie \boxed{\text{List } L}_{I_L}$$

# Merging Resources



# CoLoSL Principles



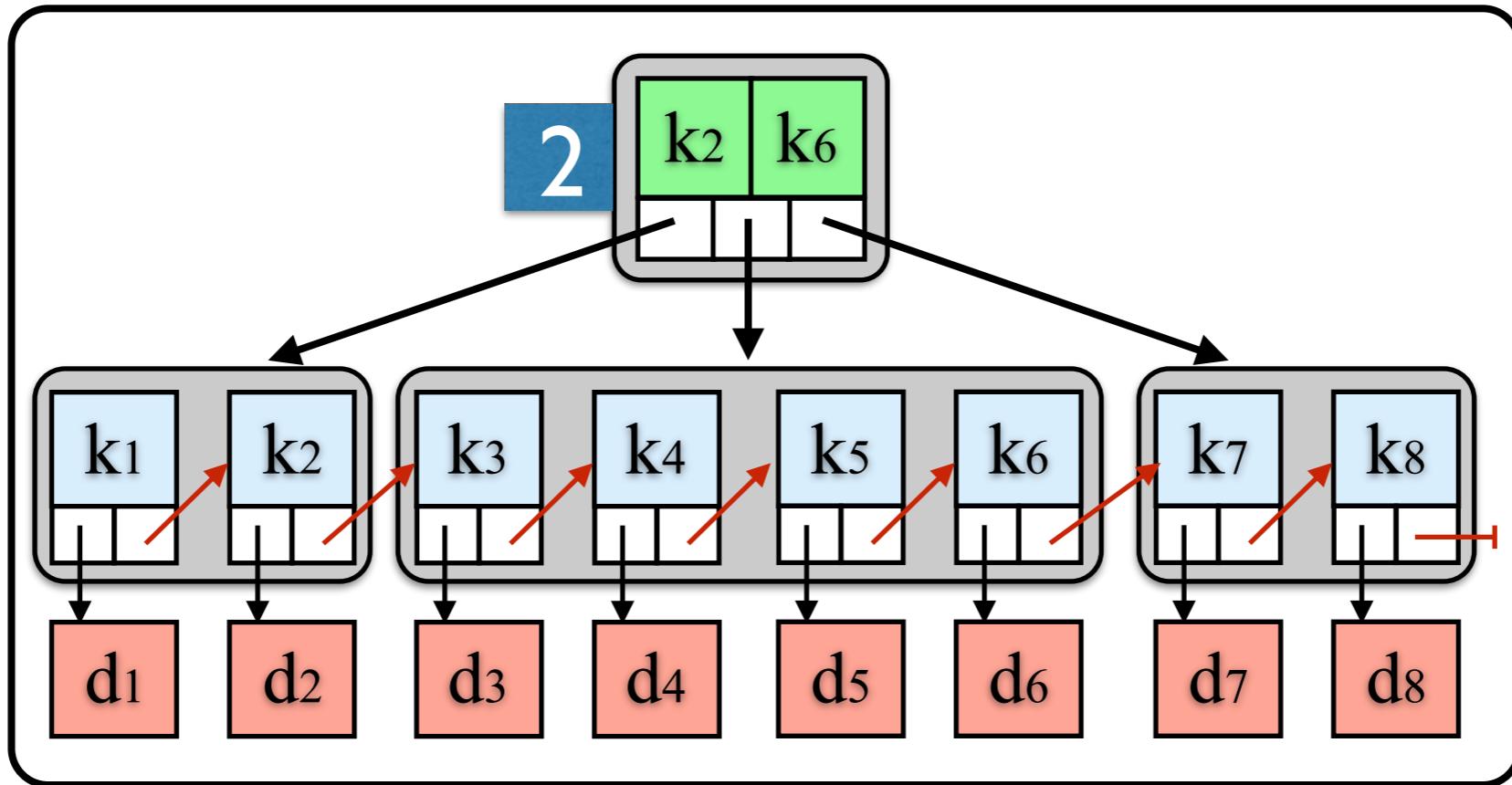
(Merge)



# Examples

- ✿ **B+ Tree**
  - ◆ Module composition; proof modularity; better abstraction
  - ◆ Proof reuse - local proofs done in the largest context possible

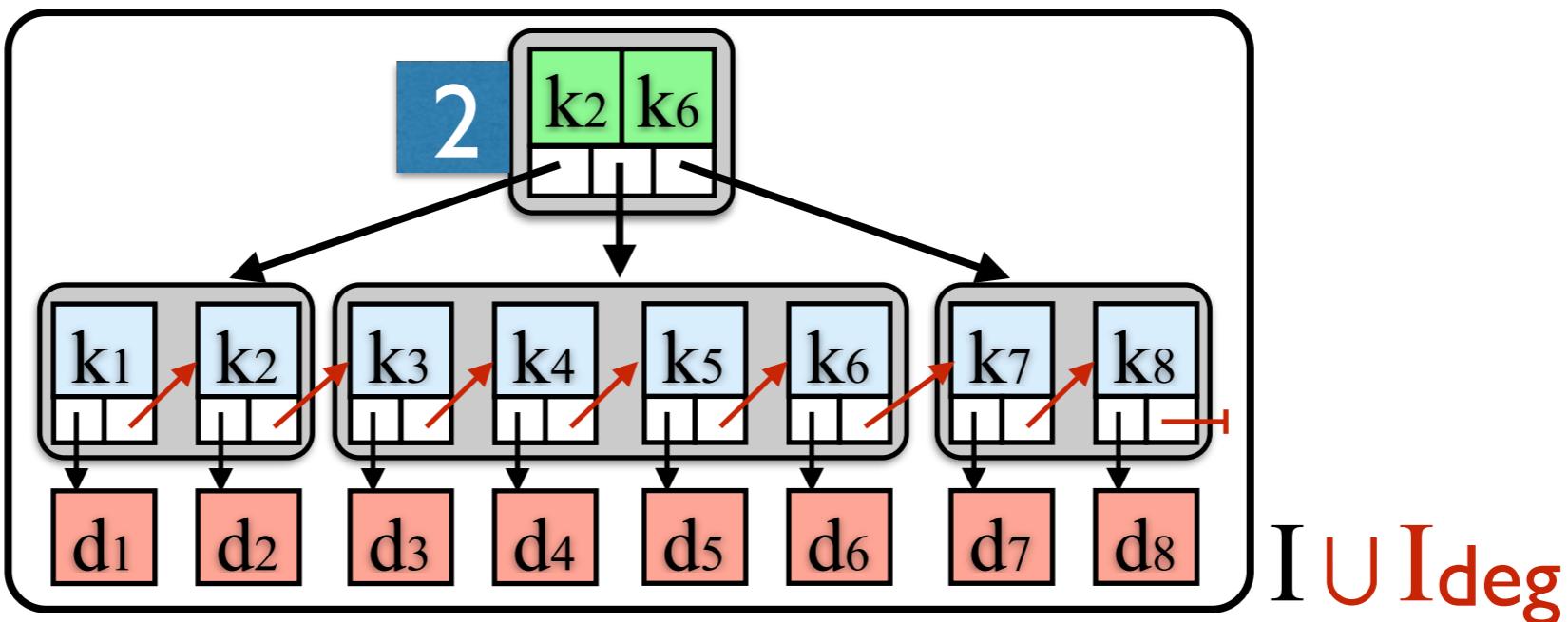
# B+ Tree: Possible Extension



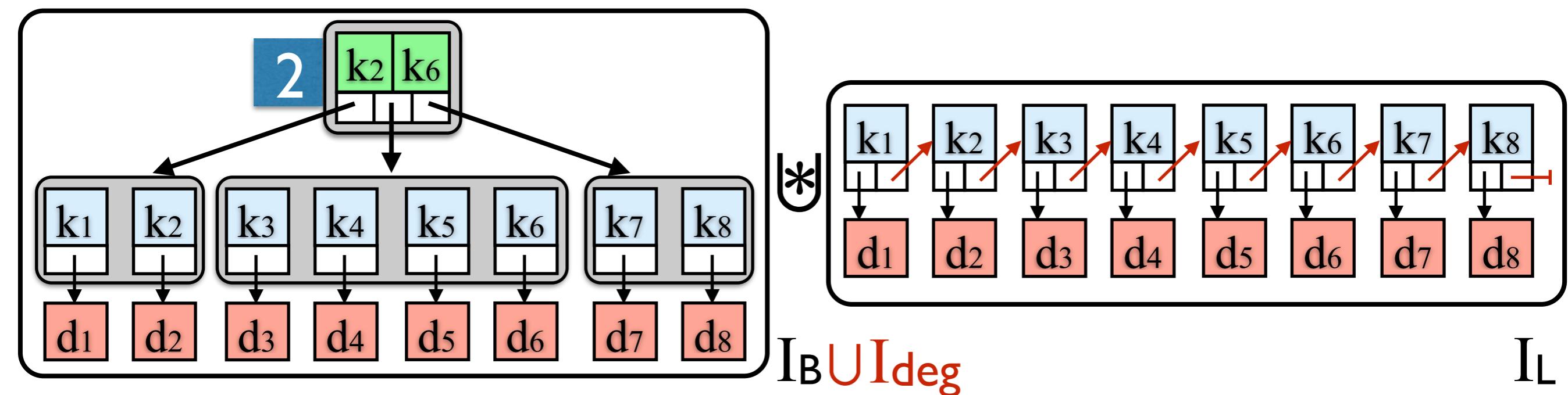
I  $\cup$  I<sub>deg</sub>

- ✿ Re-degree the tree
  - ✿ periodically change the degree ( $d$ ) for better search time
- ✿ Re-degreeing ONLY affects the tree structure (e.g.  $2 \rightsquigarrow 3$ )
- ✿ Re-degreeing does NOT affect the list structure
  - ✿ Should not have to reprove list operations

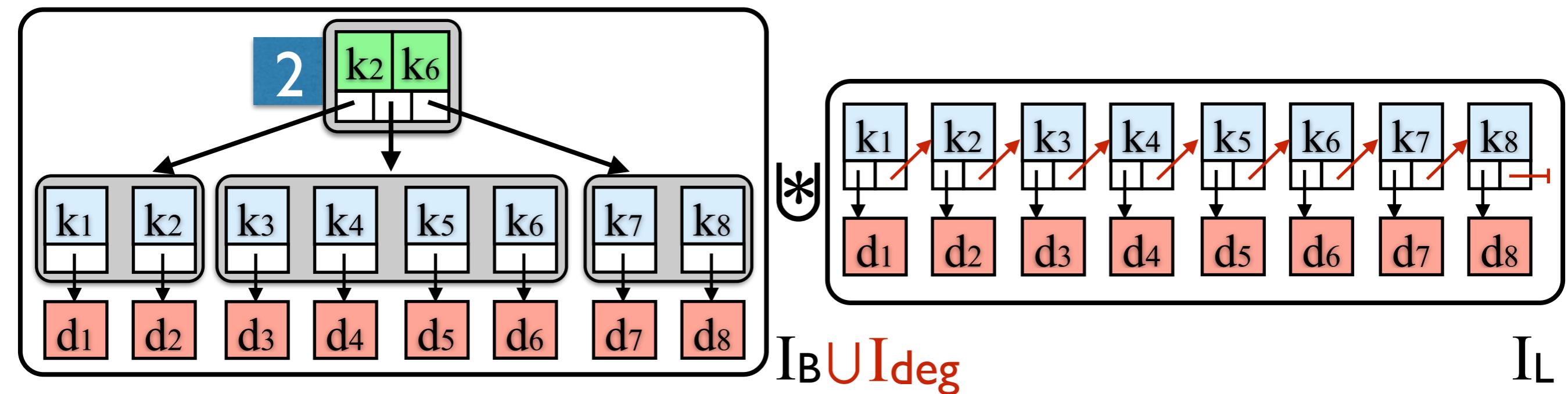
# B+ Tree: Possible Extension (CoLoSL)



$\leftrightarrow$



# B+ Tree: Possible Extension (CoLoSL)



- ❖ Re-degreeing extension only affects the tree
- ❖ Re-degreeing extension does NOT affect the list proofs
  - ❖ Can reuse the proofs as before

# Examples

- ✿ B+ Tree

- ✿ Module composition; proof modularity; better abstraction
- ✿ Proof reuse

- ✿ Concurrent List

- ✿ Dynamic extension

- ✿ Spanning Tree

- ✿ Recursive overlapping graph predicate
- ✿ Local proof; proof structures matches the algorithm

- ✿ Dijkstra's Self-stabilising Token Ring

- ✿ Local proof; proof Reuse

# Conclusions

- ✿ From OG/Rg to CAP/TaDA
  - ◆ Huge steps towards compositionality/locality
  - ◆ Are we there yet? **No!**
- ✿ CoLoSL
  - ◆ Subjective/overlapping views
  - ◆ Flexible framing on shared resource/interference
  - ◆ Dynamic extension
  - ◆ More modular; better proof reuse
  - ◆ Are we there yet? **Still No!**
    - Abstraction layers; abstract atomicity, ...

Thank you for listening!