

CoLoSL

Concurrent Local Subjective Logic

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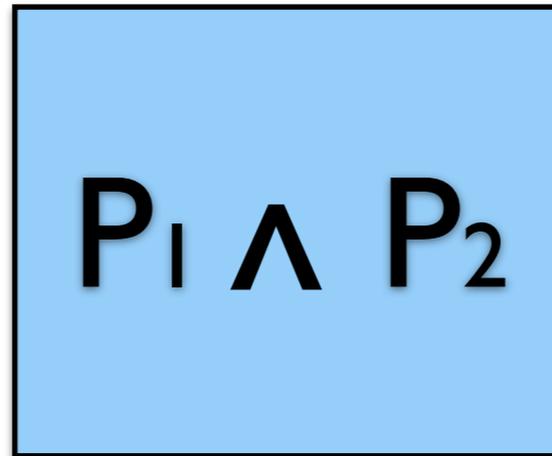
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10 April 2015

Global Shared Resources



$$\frac{\boxed{P1} \ C1 \ \boxed{Q1} \quad \boxed{P2} \ C2 \ \boxed{Q2}}{\boxed{P1 \ \wedge \ P2} \ C1 \ \parallel \ C2 \ \boxed{Q1 \ \wedge \ Q2}}$$

Global Shared Resources

$$\frac{\boxed{P1} \ C1 \ \boxed{Q1} \quad \boxed{P2} \ C2 \ \boxed{Q2}}{\boxed{P1 \wedge P2} \ C1 \parallel C2 \ \boxed{Q1 \wedge Q2}}$$

- ❖ Shared Resources

- ✦ No framing : reasoning on GLOBAL resources

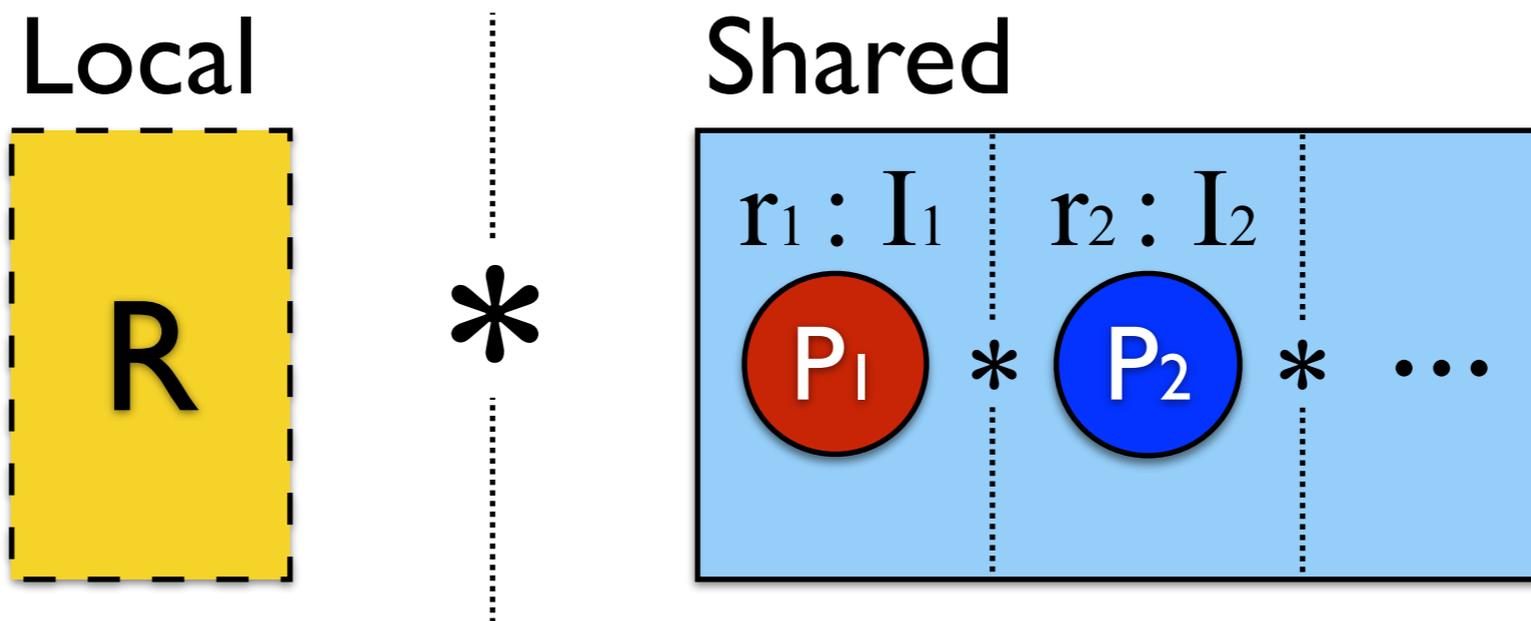
- ❖ Interference

- ✦ No framing : interference on ALL resources considered

- ❖ Extension

- ✦ No extension : cannot dynamically share resources/extend interference

Disjoint Shared Resources



Disjoint Shared Resources

$$\frac{\{\boxed{P1}\} C1 \{\boxed{Q1}\} \quad \{\boxed{P2}\} C2 \{\boxed{Q2}\}}{\{\boxed{P1} * \boxed{P2}\} C1 \parallel C2 \{\boxed{Q1} * \boxed{Q2}\}}$$

❖ Shared resources / Interference

✦ Limited framing:

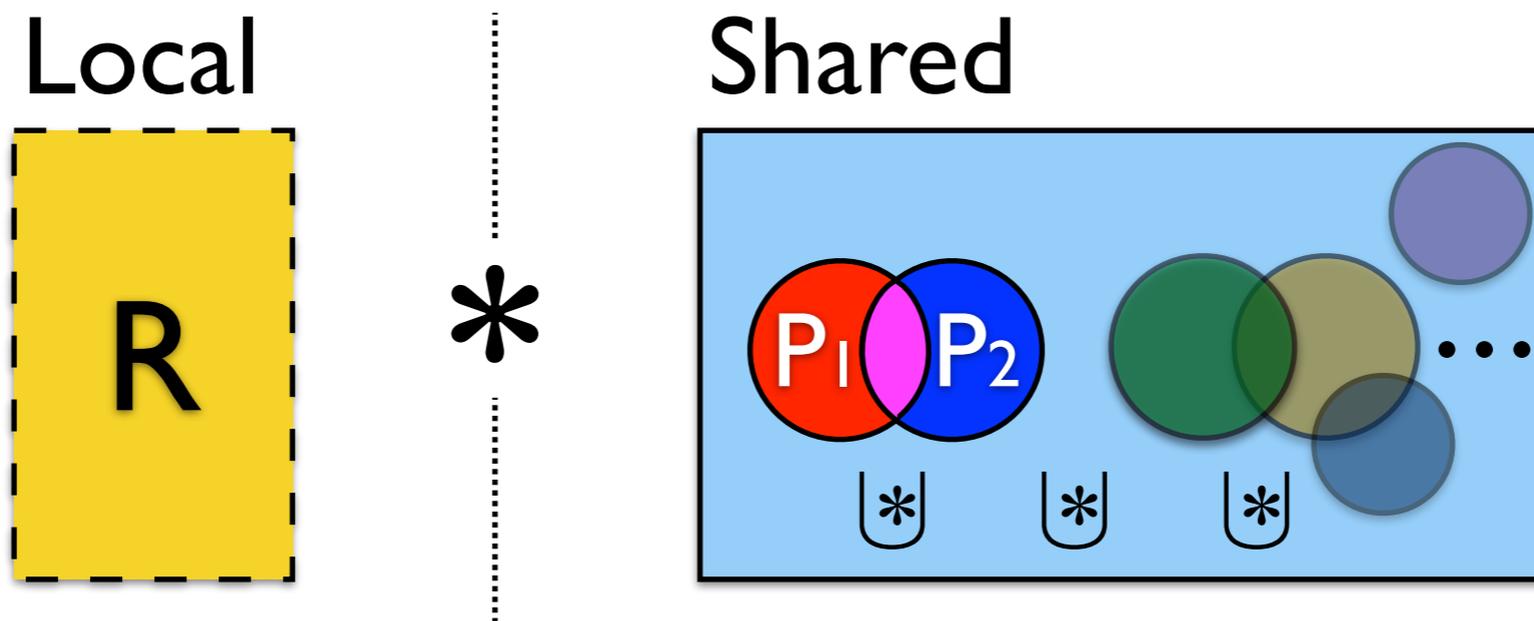
- **Static** (pre-determined) frames (regions/ invariants)
- **Physically Disjoint** frames

❖ Extension

✦ Limited extension:

- Can create new regions/ invariants
- **Cannot extend** regions with more resources/invariants

CoLoSL: Concurrent Local Subjective Logic



CoLoSL

CoLoSL: Concurrent Local Subjective Logic

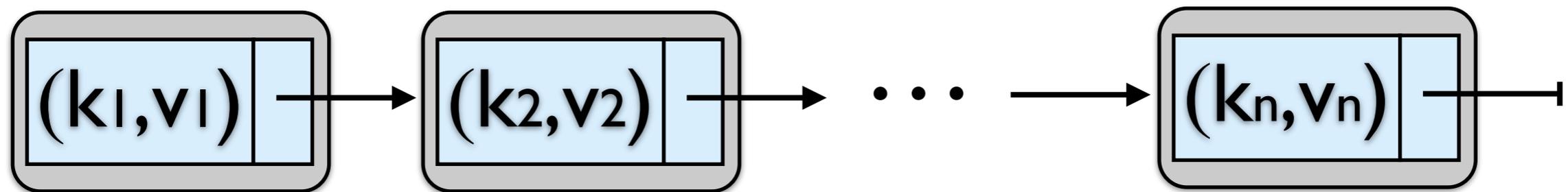
- ❖ Shared resources
 - ✦ Dynamic framing
 - ✦ Overlapping frames
- ❖ Interference
 - ✦ Dynamic framing/rewriting of interference
- ❖ Extension
 - ✦ Dynamic extension of shared state with new resource/interference

Why CoLoSL?

- ❖ Subjective/ Overlapping Shared Resources
 - ◆ Proof modularity; better abstraction

Ordered Singly Linked-List $\langle K, V \rangle$

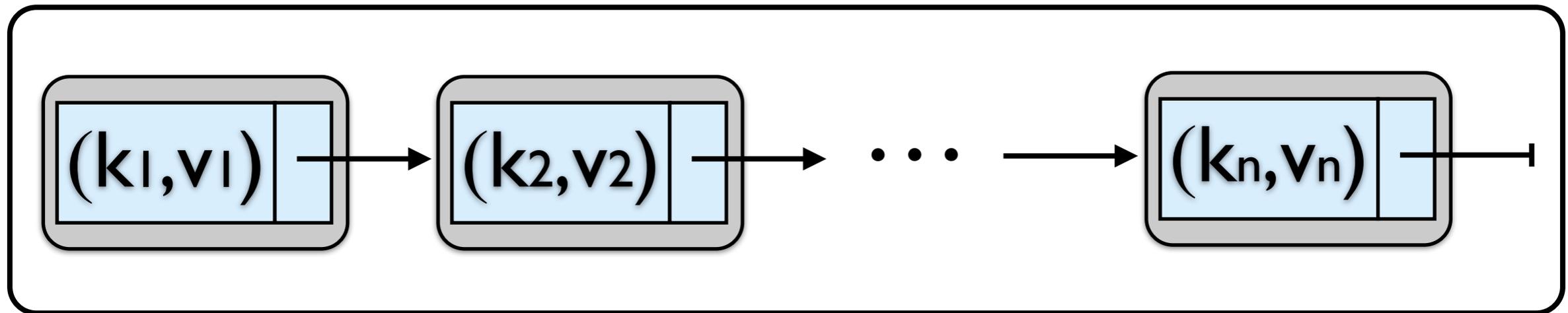
List([(k₁, v₁), (k₂, v₂), ..., (k_n, v_n)]) =



❖ insertion/removal involves pointer surgery

Concurrent Ordered Singly Linked-List<K,V>

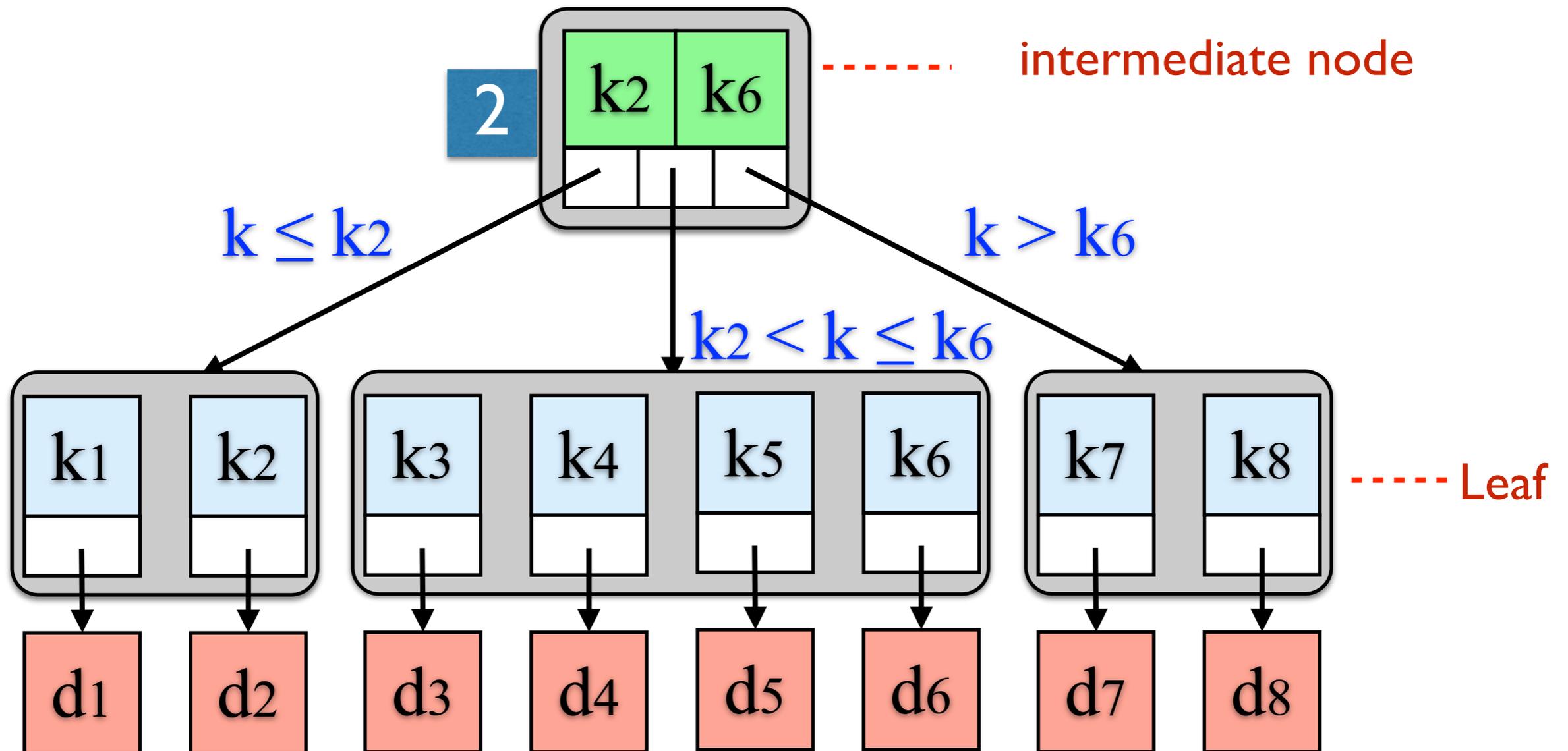
List([(k₁, v₁), (k₂, v₂), ..., (k_n, v_n)] =



$$I_L = I_{\text{add}} \cup I_{\text{rem}} \cup I_{\text{map}}$$

- ❖ insertion/removal involves pointer surgery

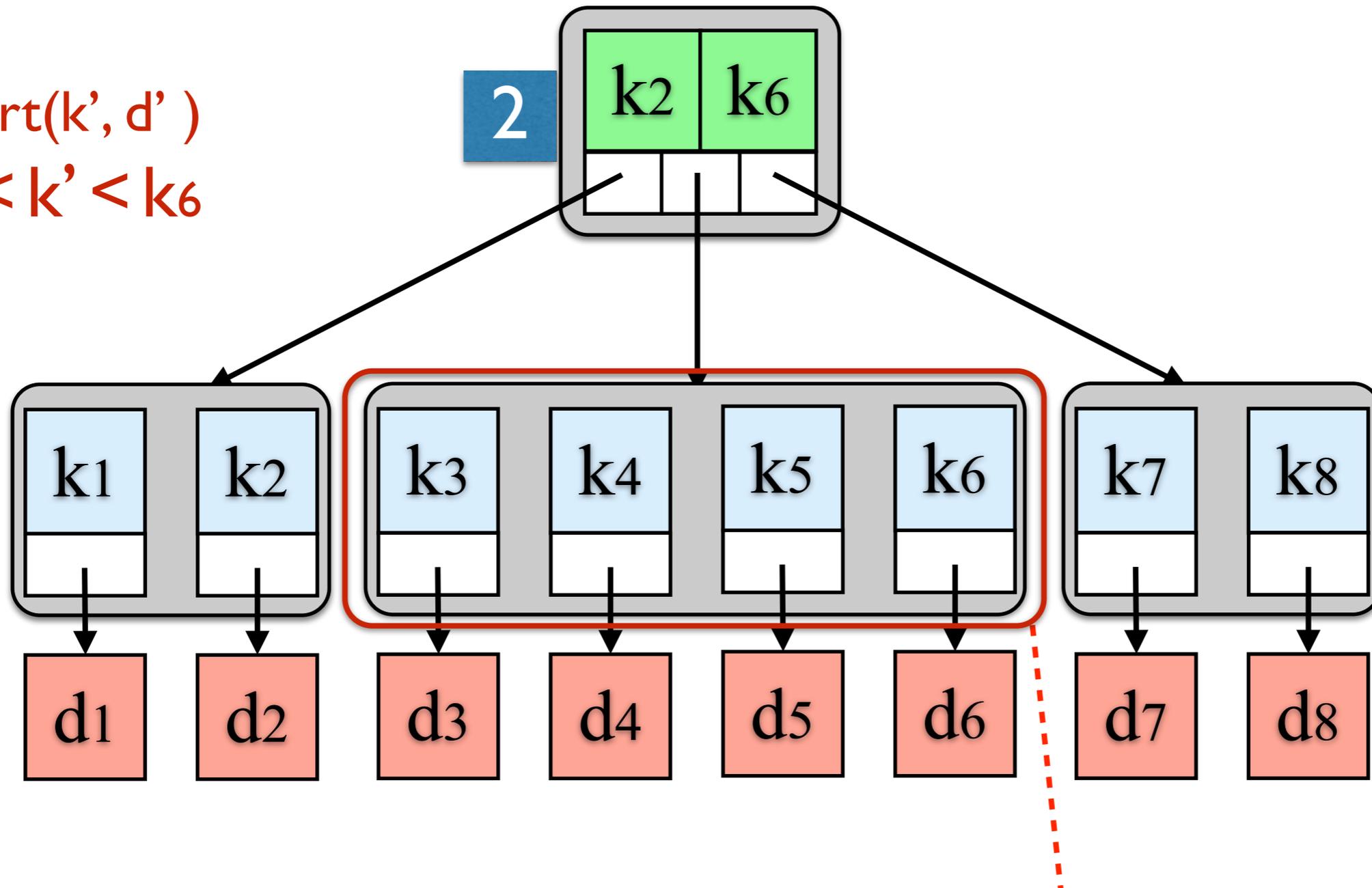
Balanced Search Tree $\langle K, V \rangle$ (Degree 2)



- ❖ **Balanced**: all immediate subtrees of a node have the same height
- ❖ **Leaf-heavy**: Data (values) stored in leaf nodes
- ❖ **Degree (d)**: no. of children (m) on each node $d \leq m \leq 2d$

Balanced Search Tree $\langle K, V \rangle$

insert(k' , d')
 $k_5 < k' < k_6$

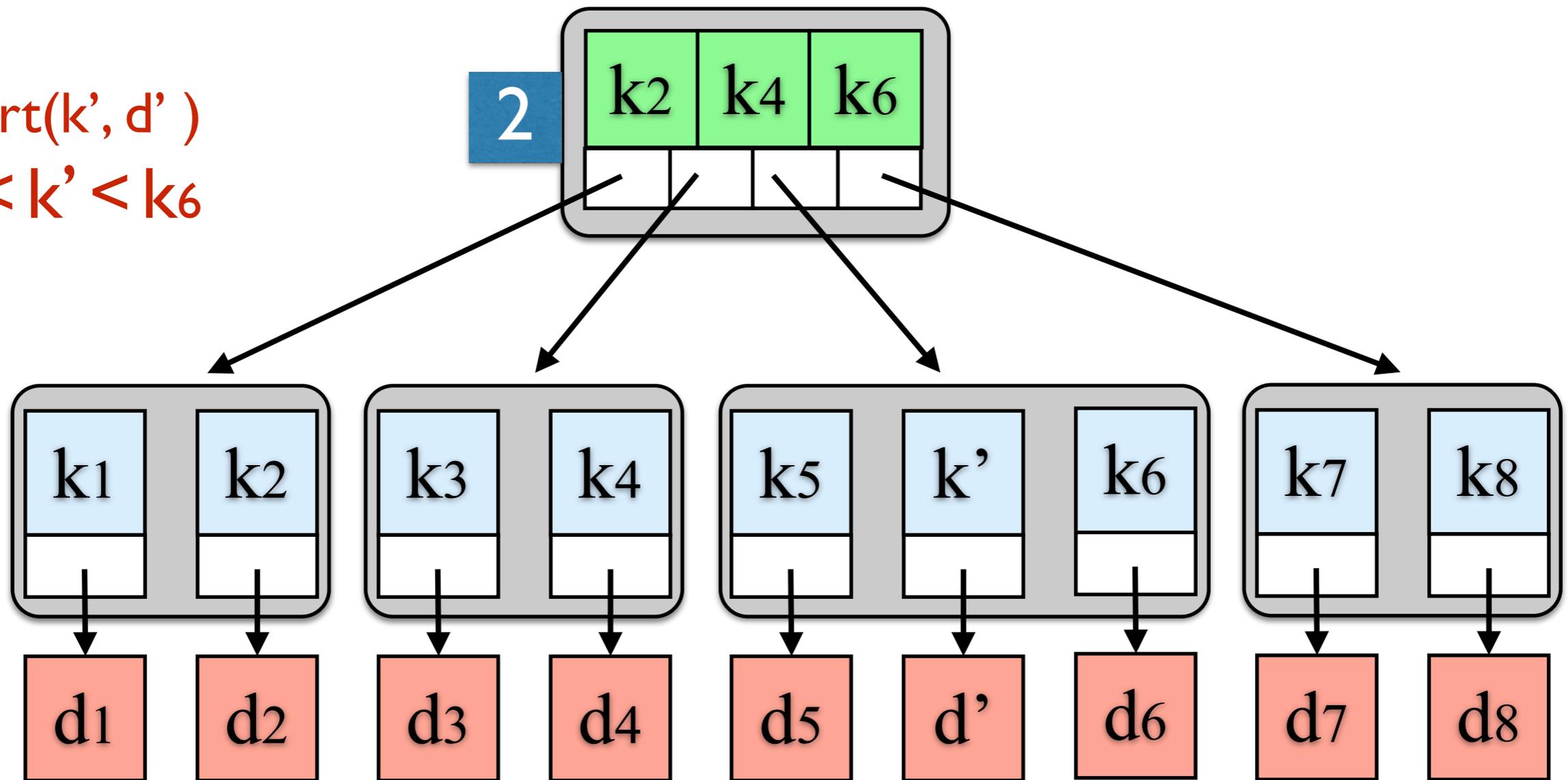


❖ Insertion may require splitting

at max capacity

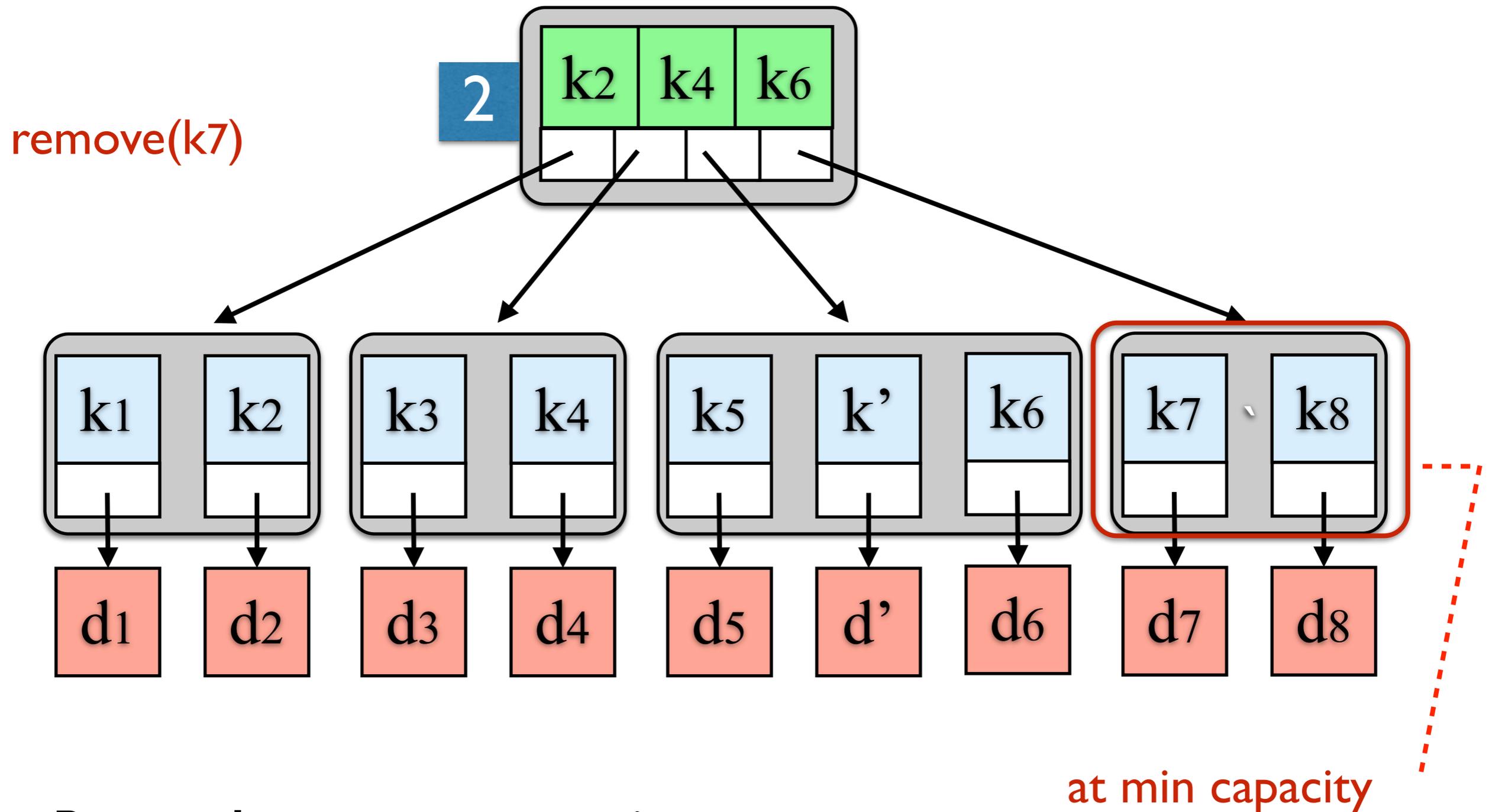
Balanced Search Tree $\langle K, V \rangle$

insert(k' , d')
 $k_5 < k' < k_6$



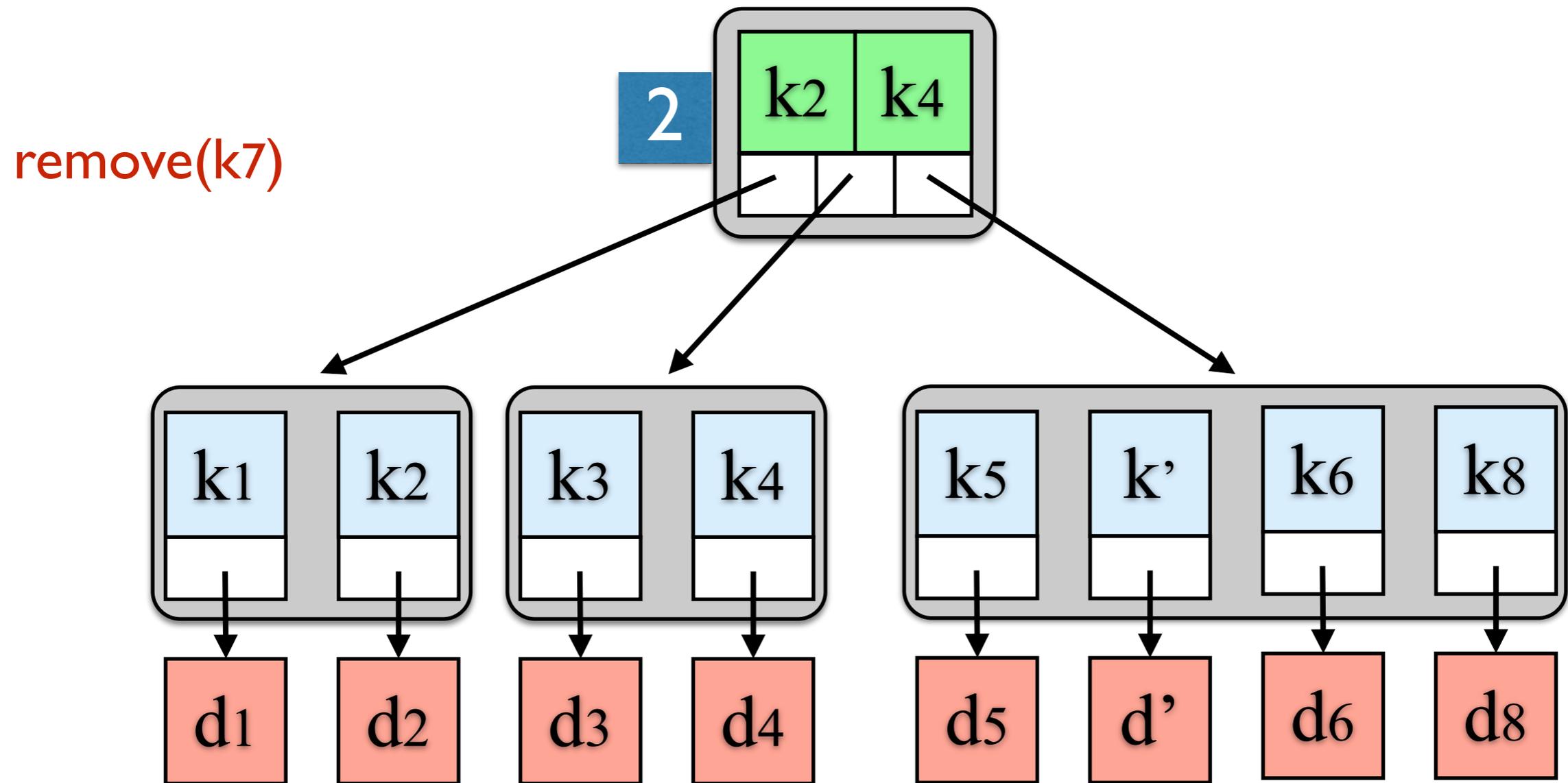
❖ Insertion may require splitting

Balanced Search Tree $\langle K, V \rangle$



❖ Removal may require merging

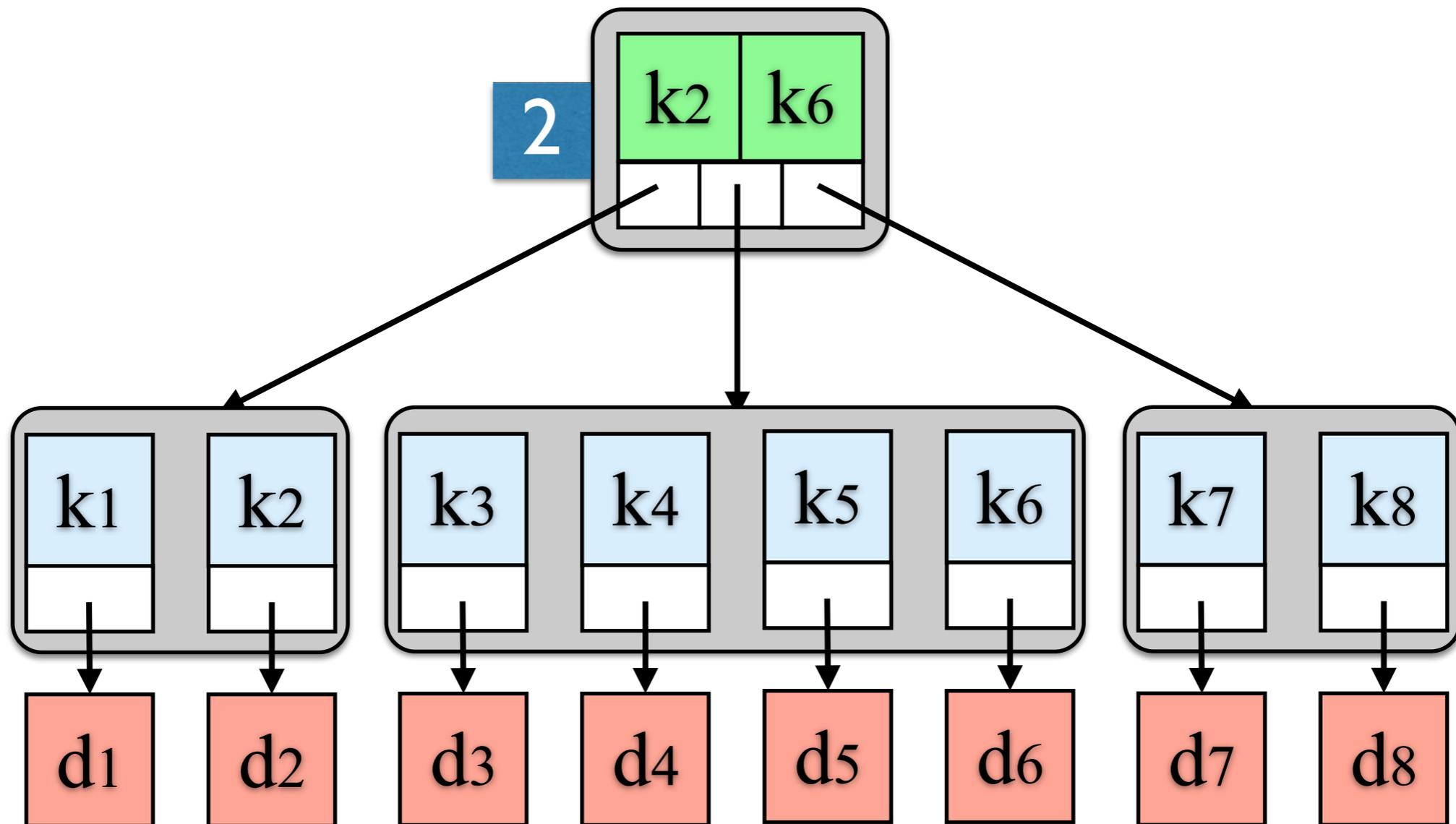
Balanced Search Tree $\langle K, V \rangle$



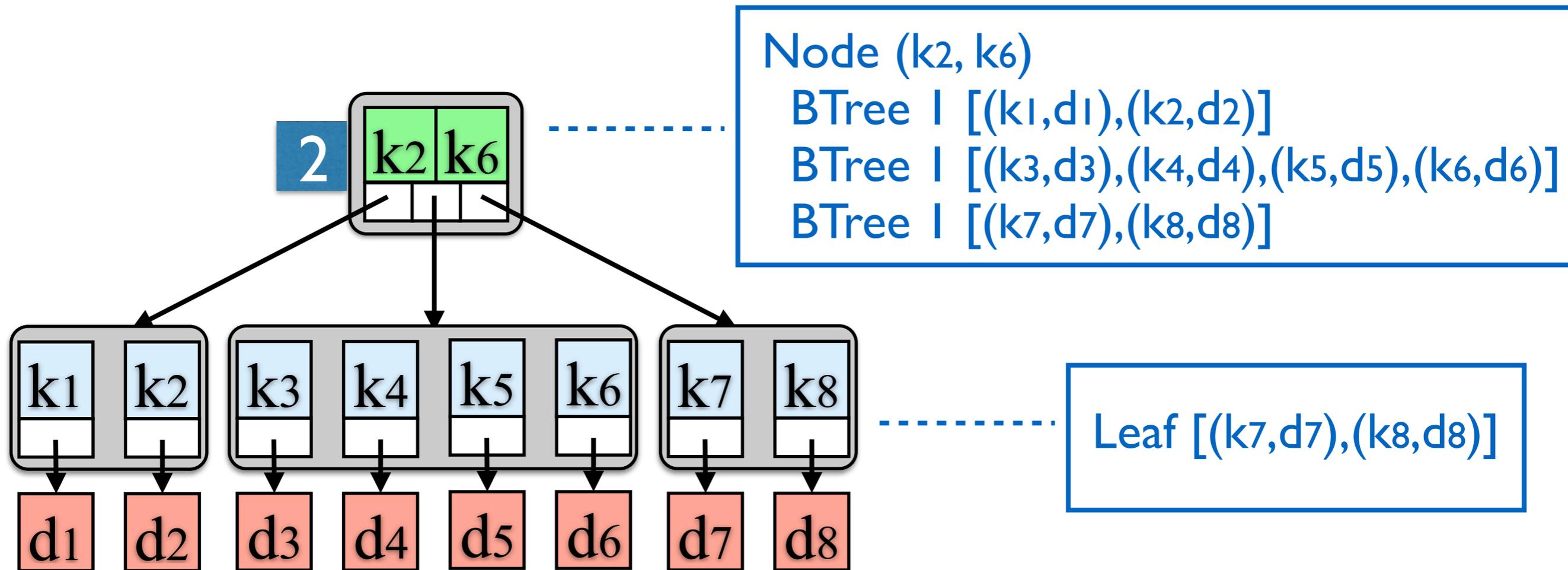
❖ Removal may require merging

Balanced Search Tree $\langle K, V \rangle$

BTree 2 $([(k_1, d_1) \dots (k_8, d_8)]) =$



Balanced Search Tree $\langle K, V \rangle$



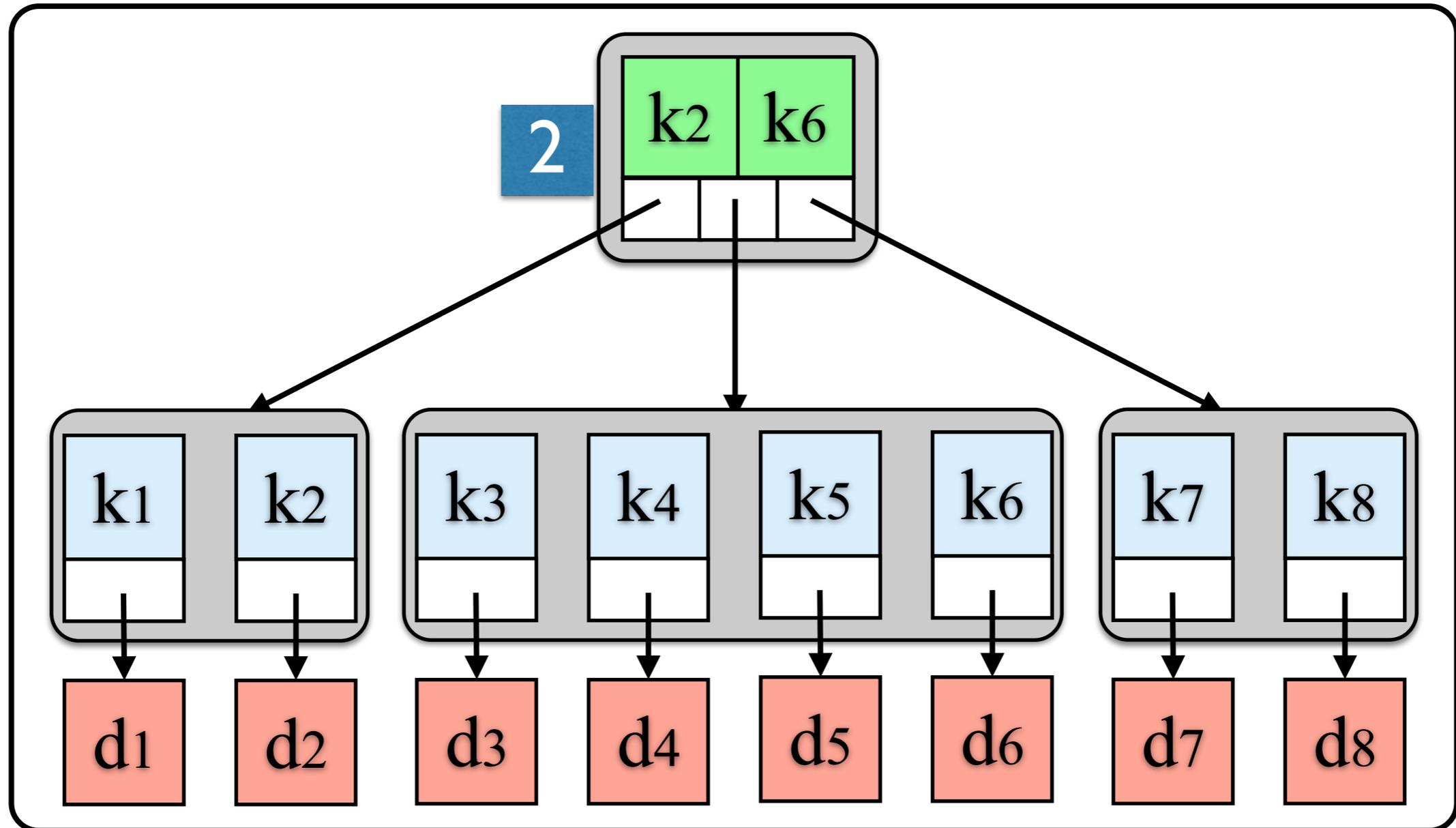
BTree (h+1) L =

$\exists k_1, \dots, k_m, L_1, \dots, L_{(m+1)} . L = L_1 \uplus \dots \uplus L_{(m+1)}$
 Node (k1, ..., km) (BTree h L1) ... (BTree h L(m+1))

BTree 1 L = Leaf L

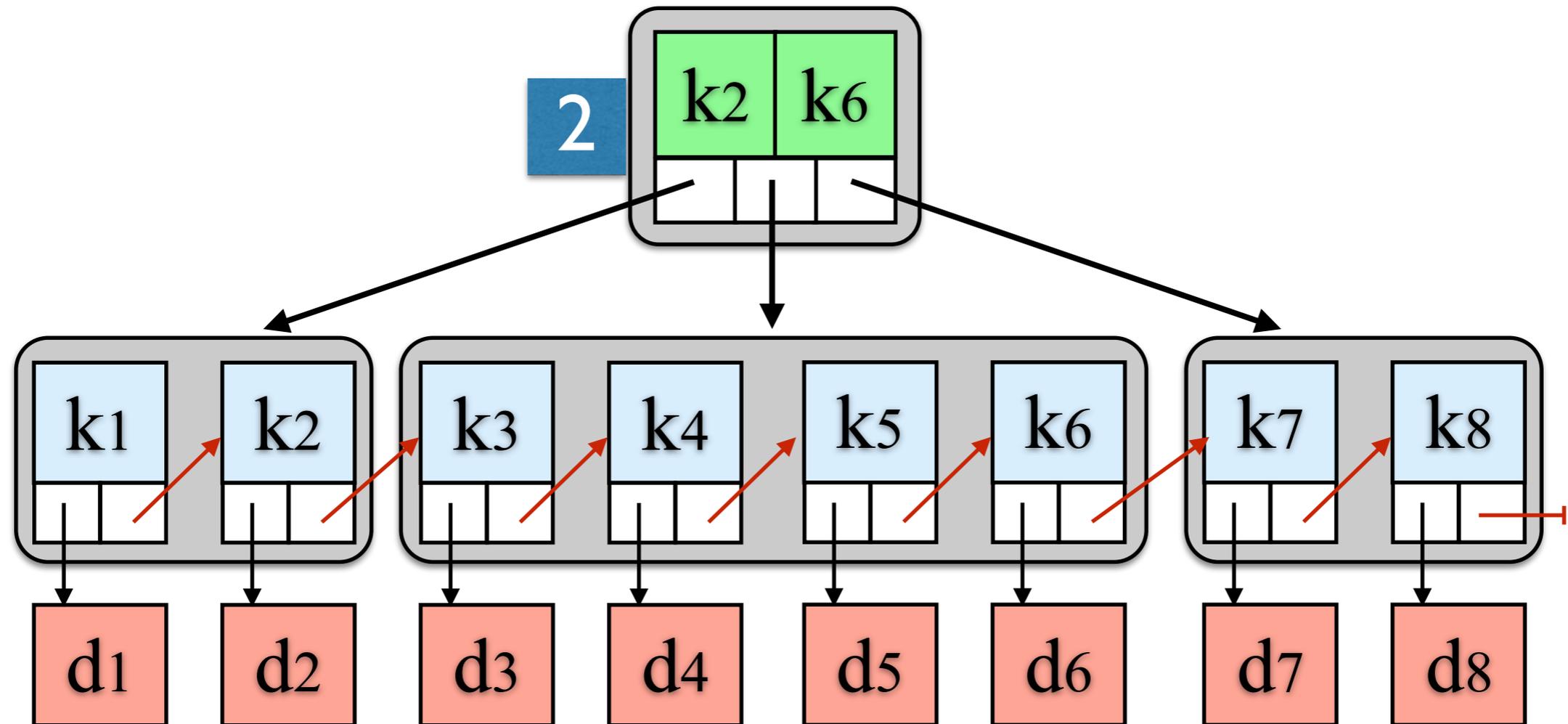
Concurrent Balanced Search Tree $\langle K, V \rangle$

BTree([(k1, d1) ... (k8, d8)]) =

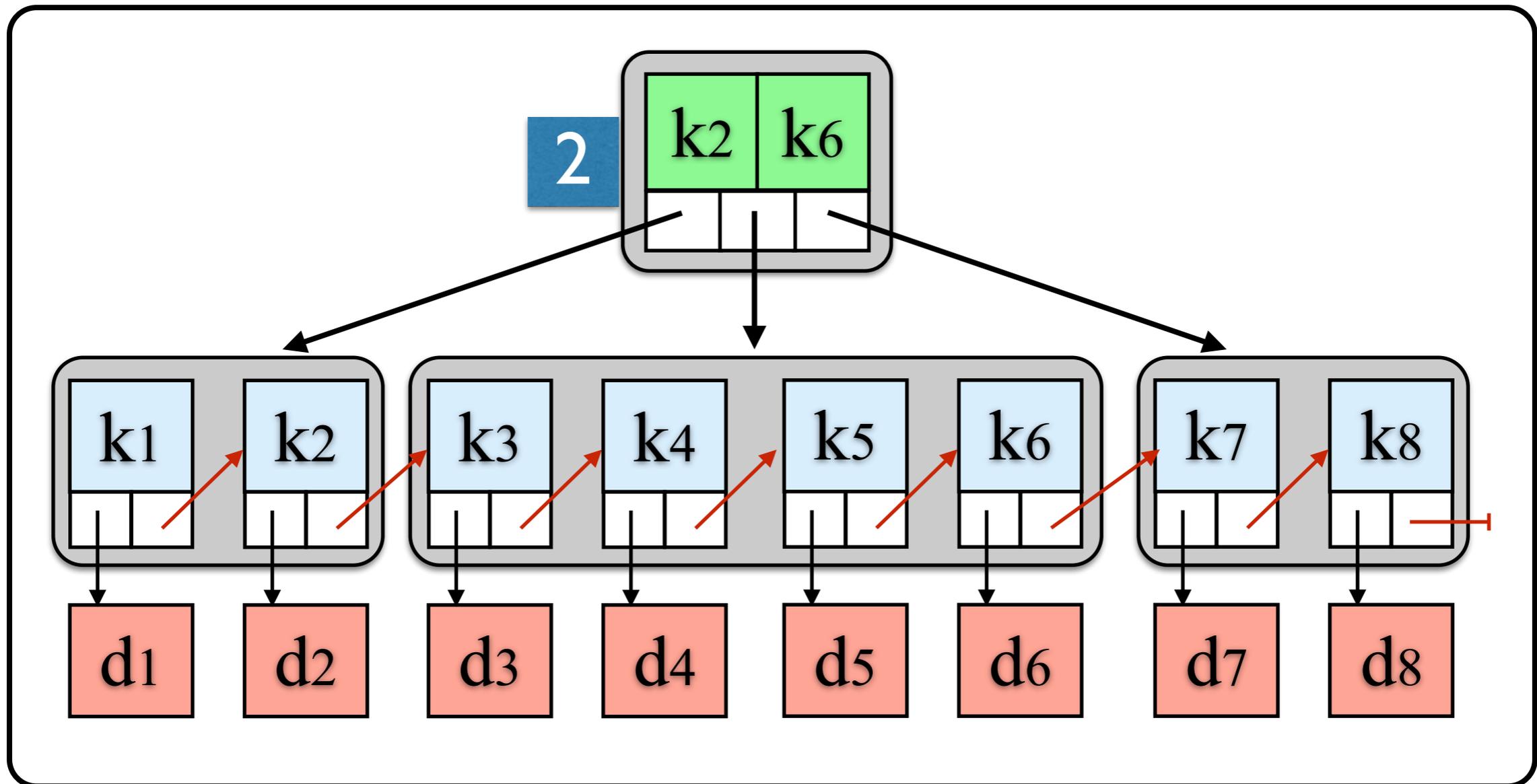


$$I_B = I_{ins} \cup I_{del} \cup I_{srch}$$

B+ Tree $\langle K, V \rangle$



Concurrent B+ Tree $\langle K, V \rangle$

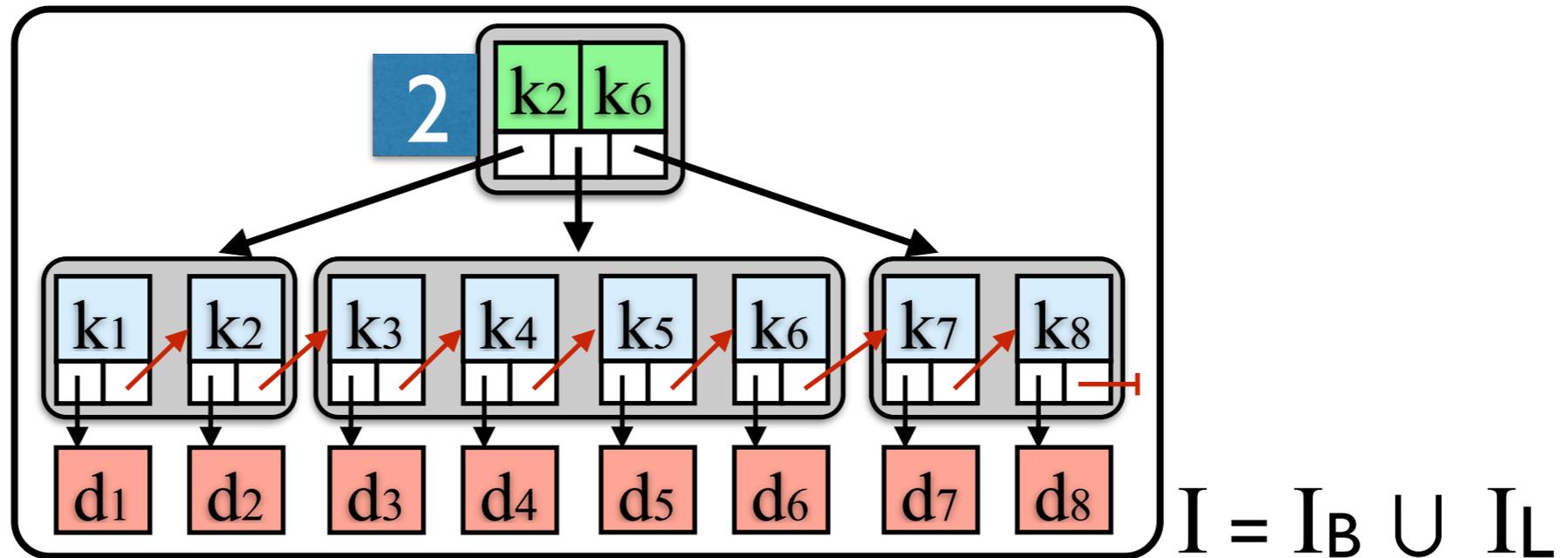


$$I = I_B \cup I_L$$

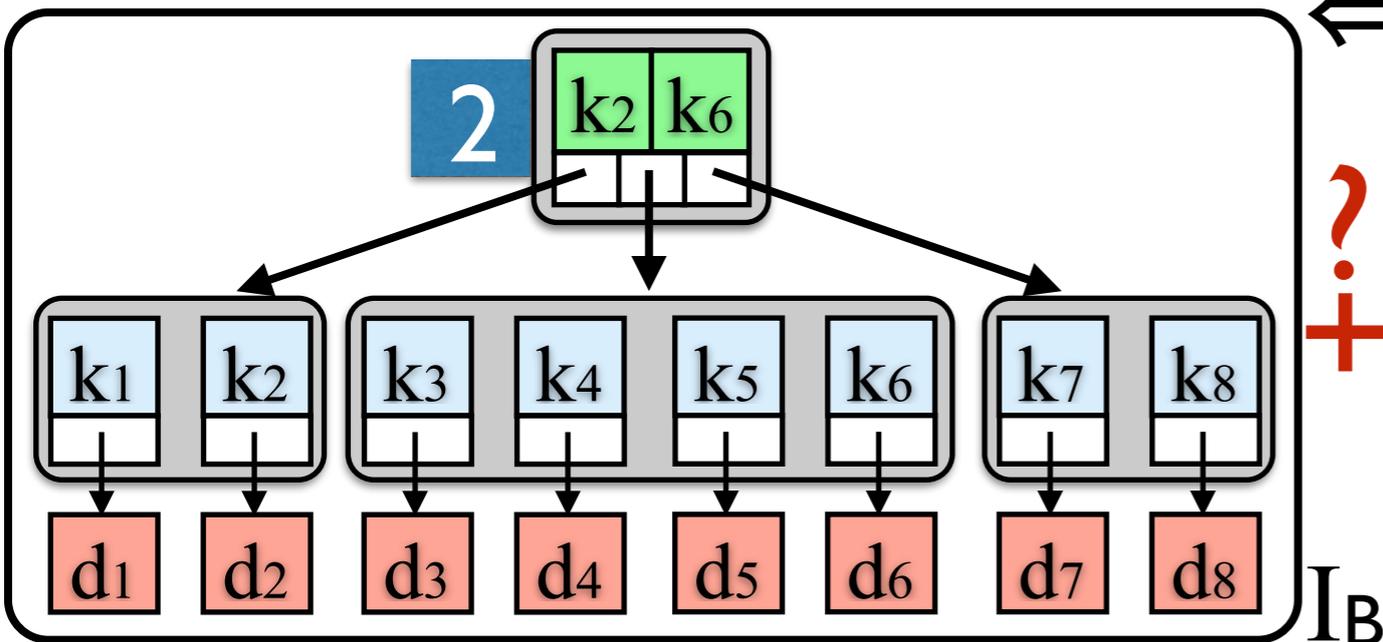
Concurrent B+ Tree $\langle K, V \rangle$ Wish List

❖ Module Composition

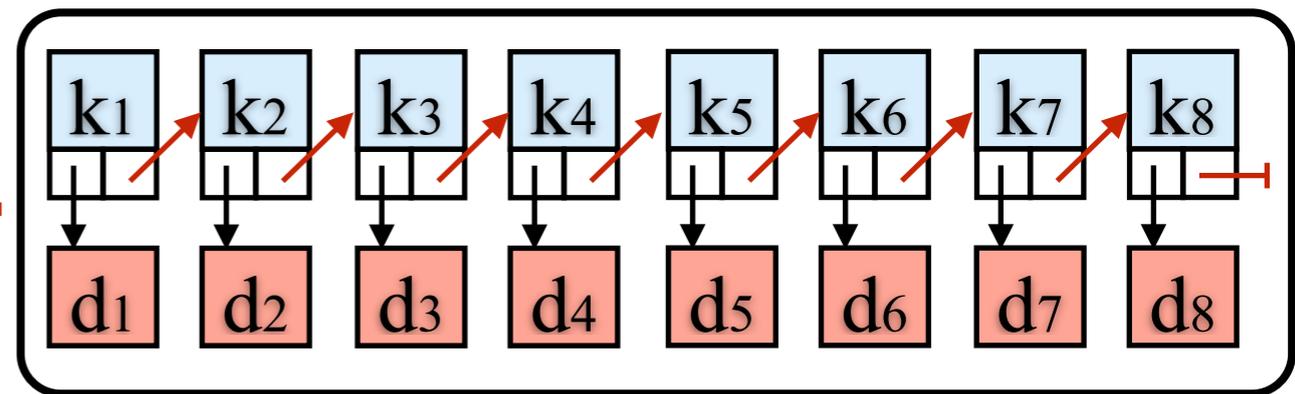
$$\boxed{\text{B+ Tree } h \text{ L}}_{I_B \cup I_L} \Leftrightarrow \boxed{\text{BTree } h \text{ L}}_{I_B} \overset{?}{+} \boxed{\text{List L}}_{I_L}$$



\Leftrightarrow



$\overset{?}{+}$

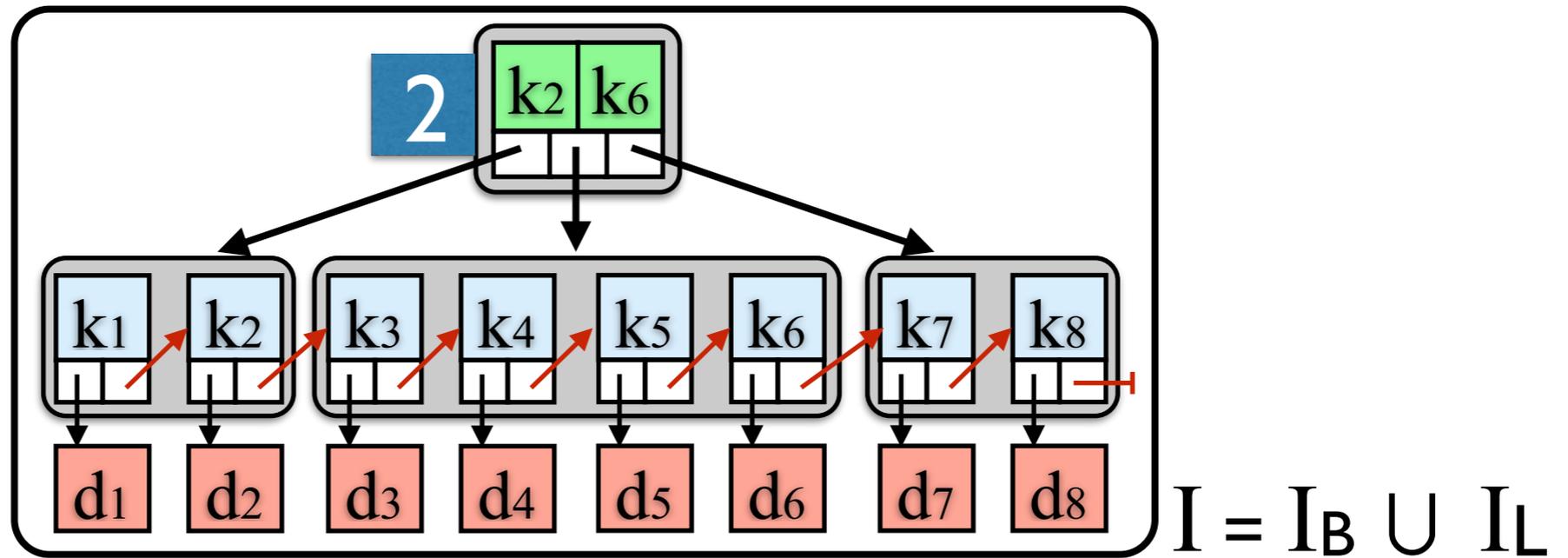


Concurrent B+ Tree $\langle K, V \rangle$

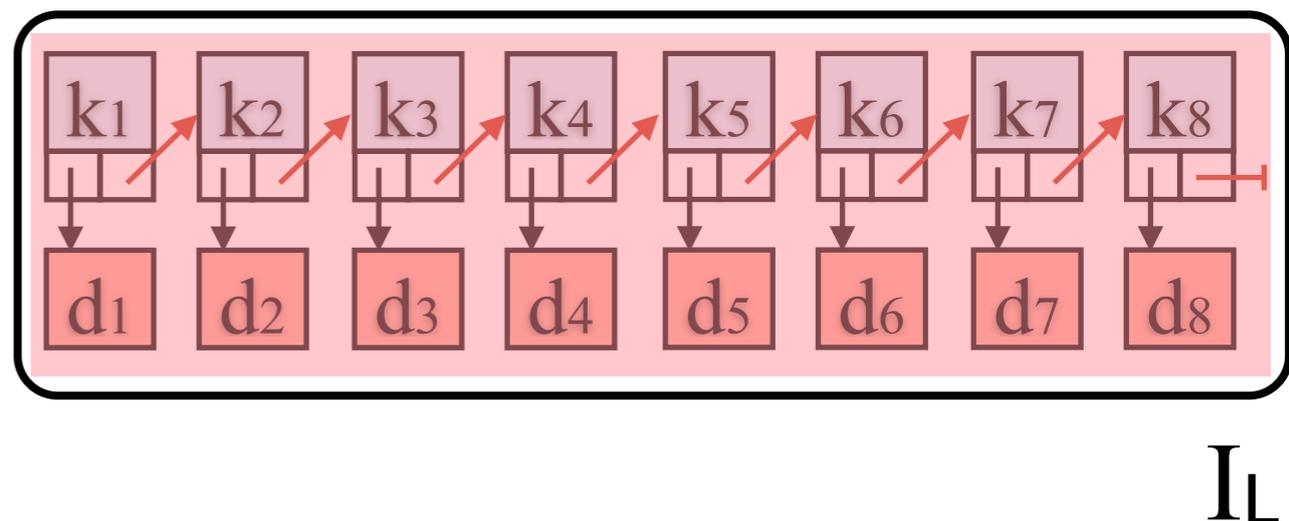
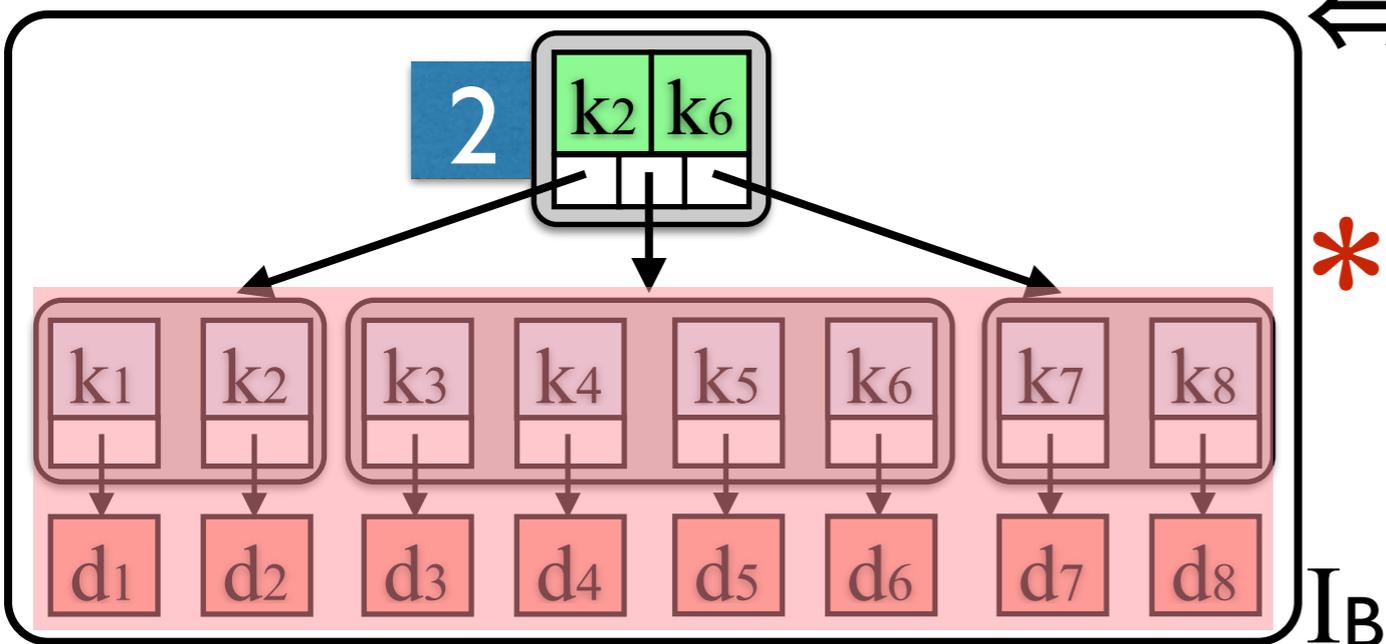
(Existing Approaches)

❖ Module Composition

$$\boxed{\text{B+ Tree } h \text{ L}}_{I_B \cup I_L} \Leftrightarrow \boxed{\text{BTree } h \text{ L}}_{I_B} * \boxed{\text{List L}}_{I_L}$$



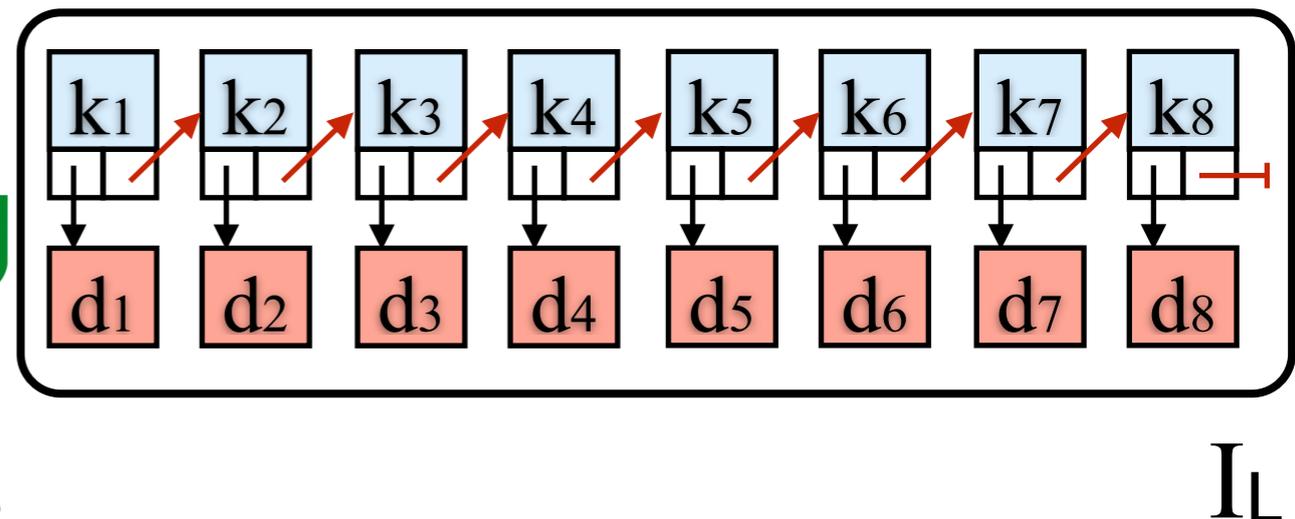
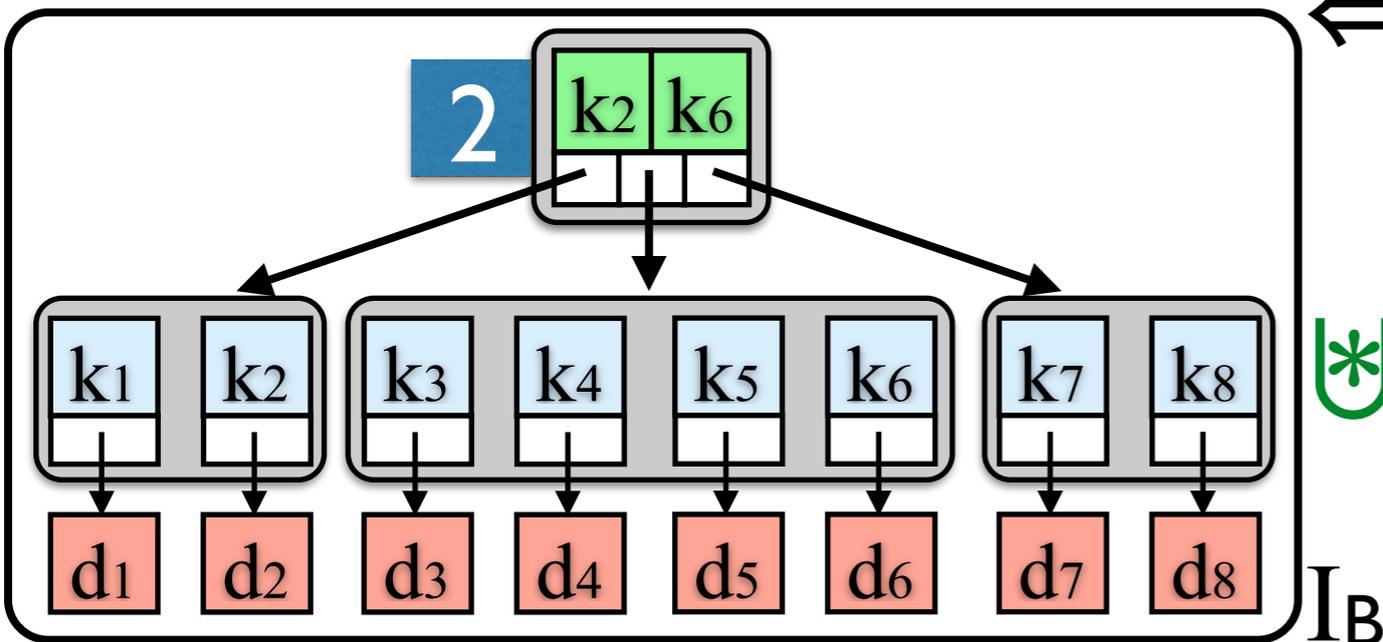
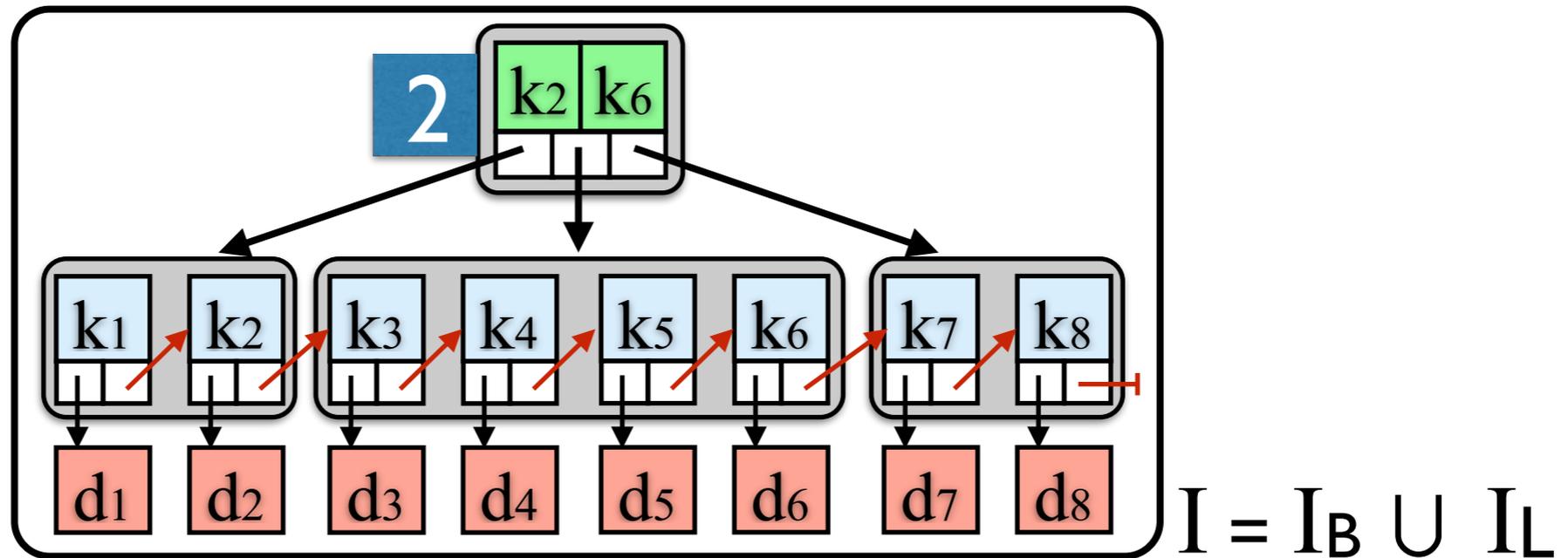
\Leftrightarrow X



Concurrent B+ Tree $\langle K, V \rangle$ (CoLoSL)

❖ Module Composition

$$\boxed{\text{B+ Tree } h \text{ L}}_{I_B \cup I_L} \Leftrightarrow \boxed{\text{BTree } h \text{ L}}_{I_B} * \boxed{\text{List L}}_{I_L}$$



B+ Tree $\langle K, V \rangle$ Wish List

❖ Module Composition

$$\boxed{\text{B+ Tree h L}}_{I_B \cup I_L} \Leftrightarrow \boxed{\text{BTree h L}}_{I_B} \ast \boxed{\text{List L}}_{I_L}$$

❖ Proof Modularity (overlapping frames)

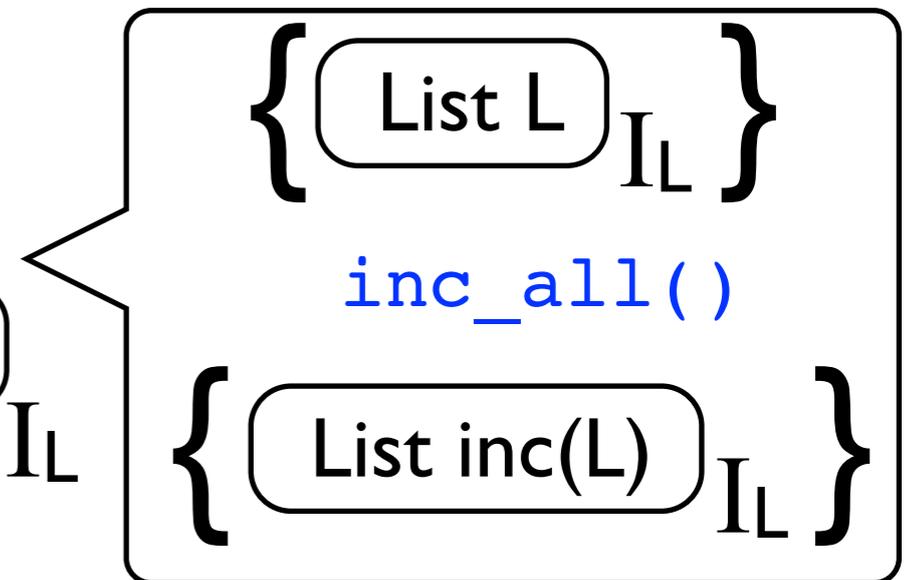
◆ Reuse list-only operation (e.g. map)

$$\boxed{\text{BTree h L}}_{I_B} \ast \boxed{\text{List L}}_{I_L} \xRightarrow{\text{(frame)}} \boxed{\text{List L}}_{I_L}$$

◆ Reuse tree-only operation (e.g. search)

$$\boxed{\text{BTree h L}}_{I_B} \ast \boxed{\text{List L}}_{I_L} \xRightarrow{\text{(frame)}} \boxed{\text{BTree h L}}_{I_B}$$

◆ Combine tree/list module operations to implement B+ tree operations (e.g. remove)



CoLoSL Principles

$$\boxed{\text{B+ Tree h L}}_{I_B \cup I_L} \Leftrightarrow \boxed{\text{BTree h L}}_{I_B} * \boxed{\text{List L}}_{I_L}$$

CoLoSL Principles

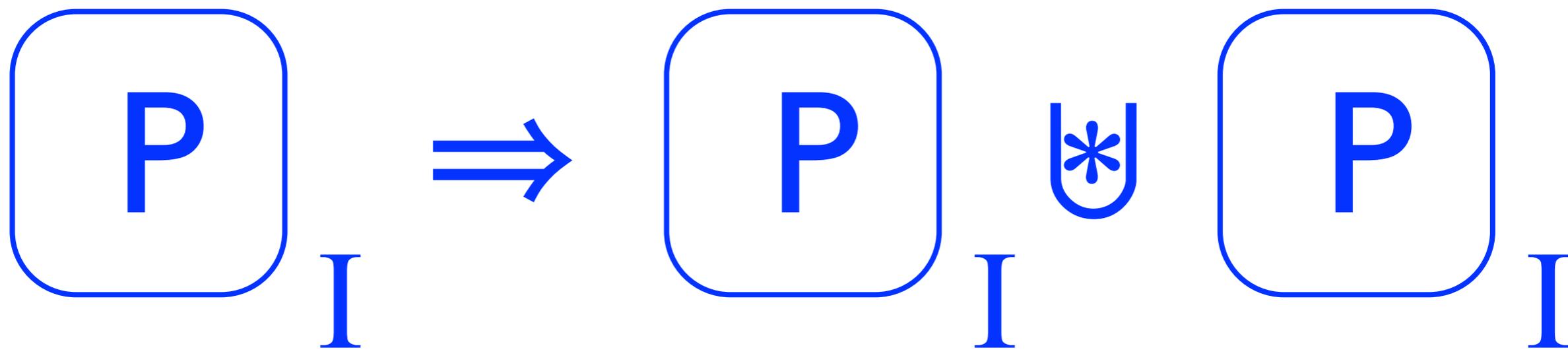
$\text{B+ Tree } h \text{ L}$ $I_{BUI}L$



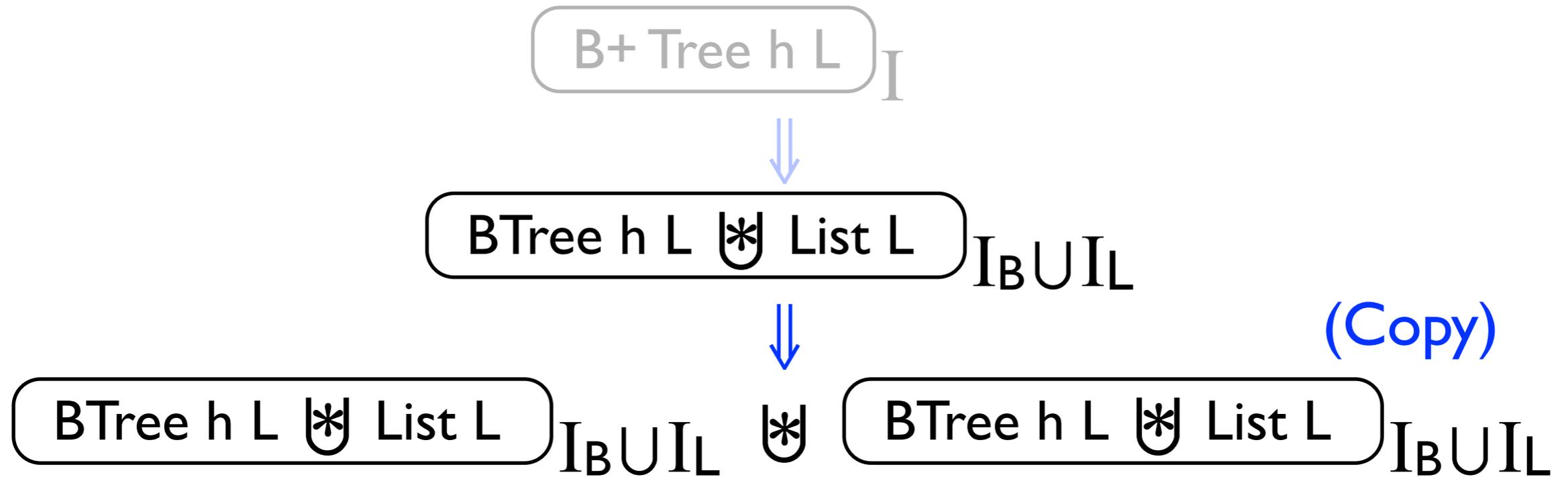
$\text{BTree } h \text{ L} \ast \text{List } L$ $I_{BUI}L$

$\text{BTree } h \text{ L}$ I_B \ast $\text{List } L$ I_L

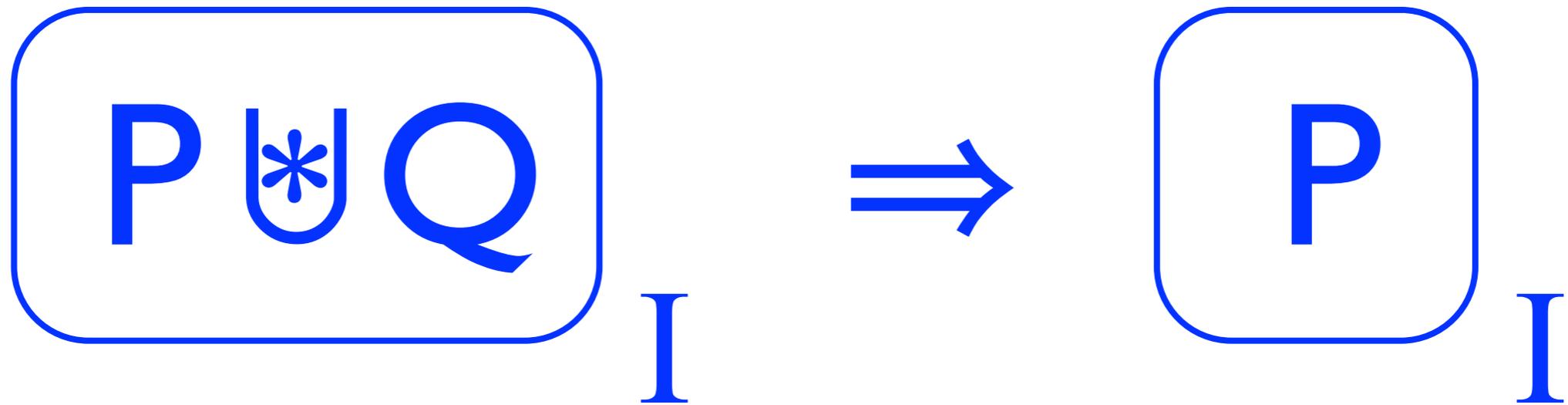
Duplicating Resources



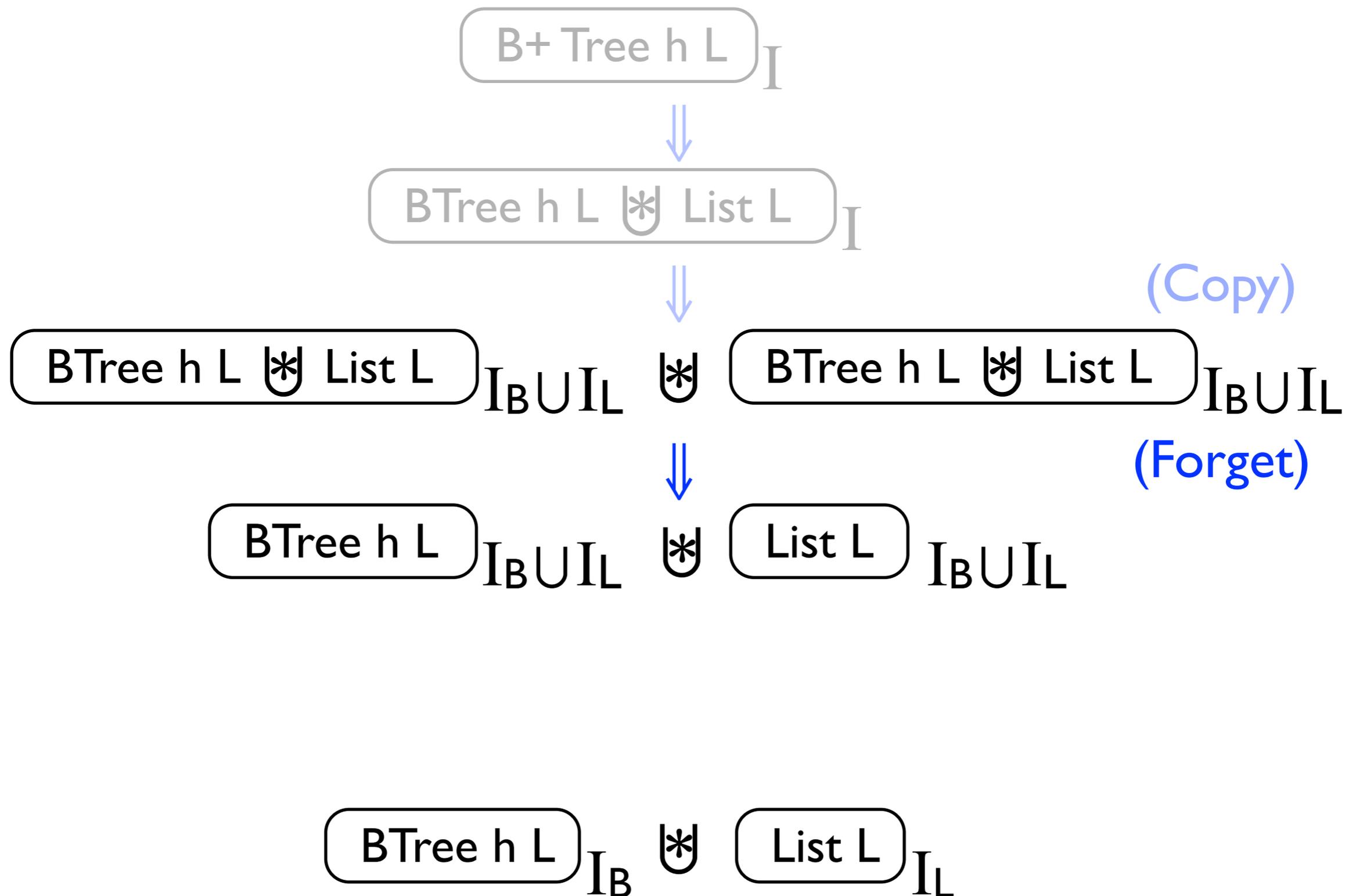
CoLoSL Principles



Forgetting Resources



CoLoSL Principles

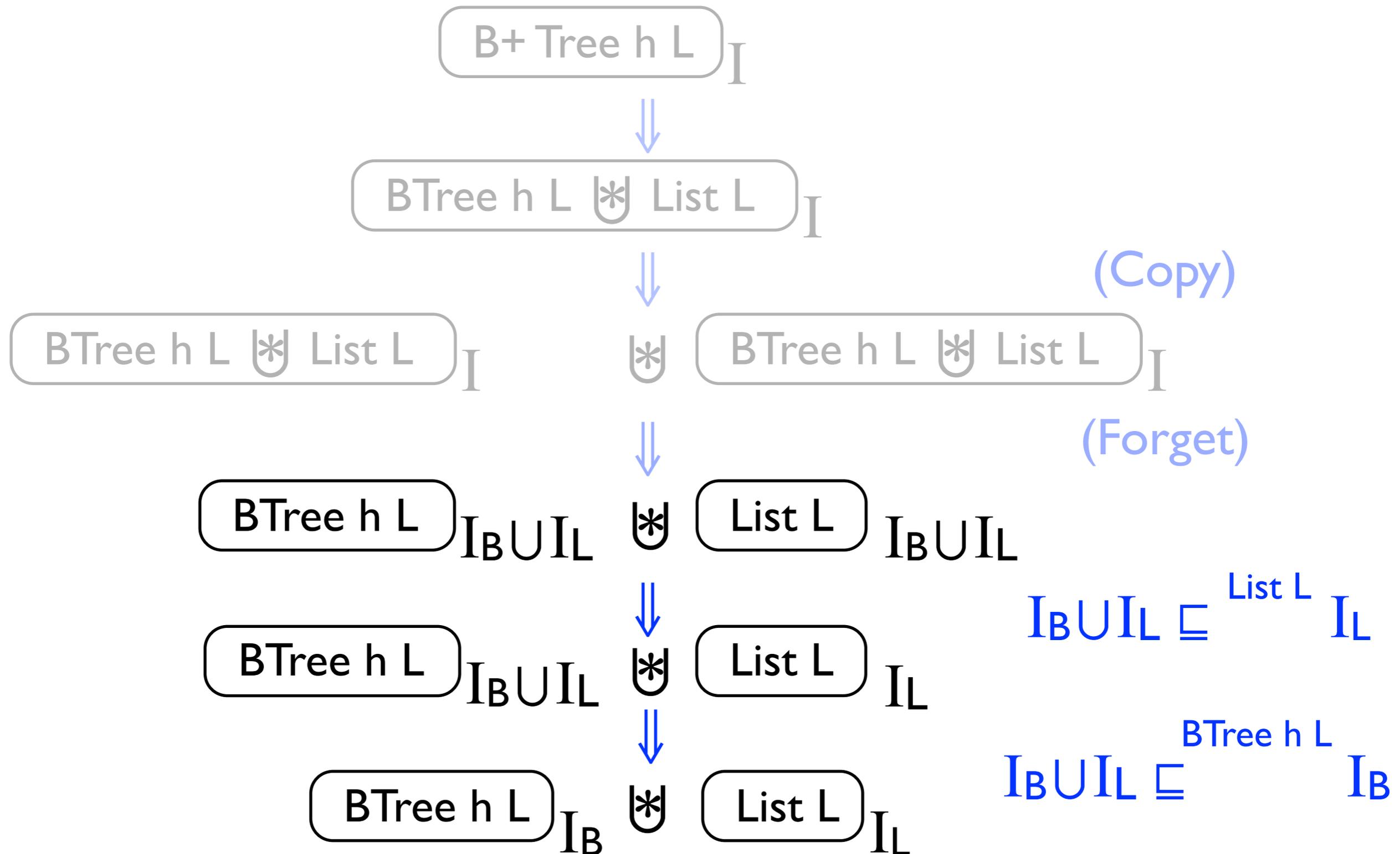


Forgetting Interference (Shift)

if $I \sqsubseteq^P I'$

then $\boxed{P}_I \Rightarrow \boxed{P}_{I'}$

CoLoSL Principles

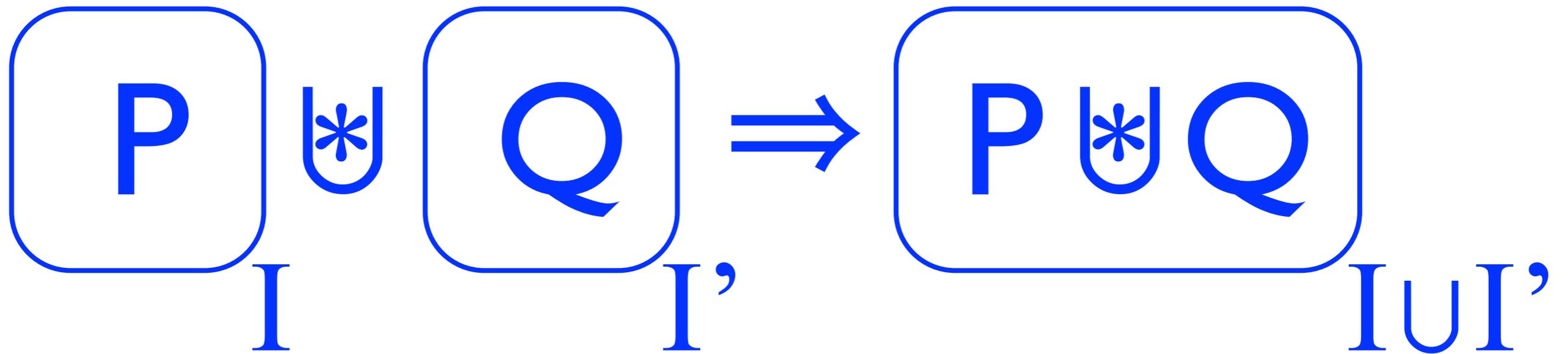


CoLoSL Principles

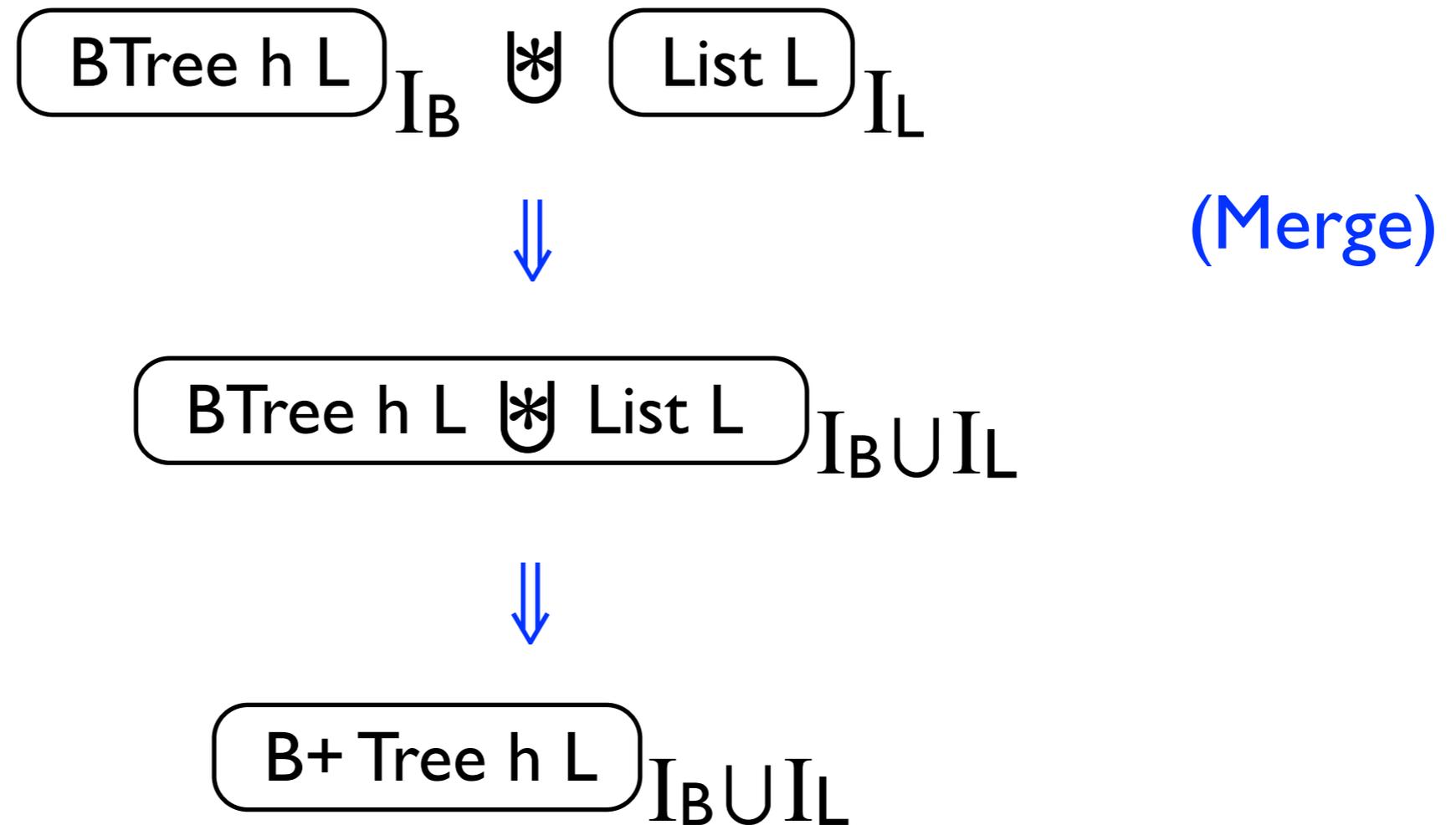
$$\boxed{\text{B+ Tree h L}}_{I_B \cup I_L} \Rightarrow \boxed{\text{BTree h L}}_{I_B} \ast \boxed{\text{List L}}_{I_L}$$

$$\boxed{\text{B+ Tree h L}}_{I_B \cup I_L} \Leftarrow \boxed{\text{BTree h L}}_{I_B} \ast \boxed{\text{List L}}_{I_L}$$

Merging Resources



CoLoSL Principles



Why CoLoSL?

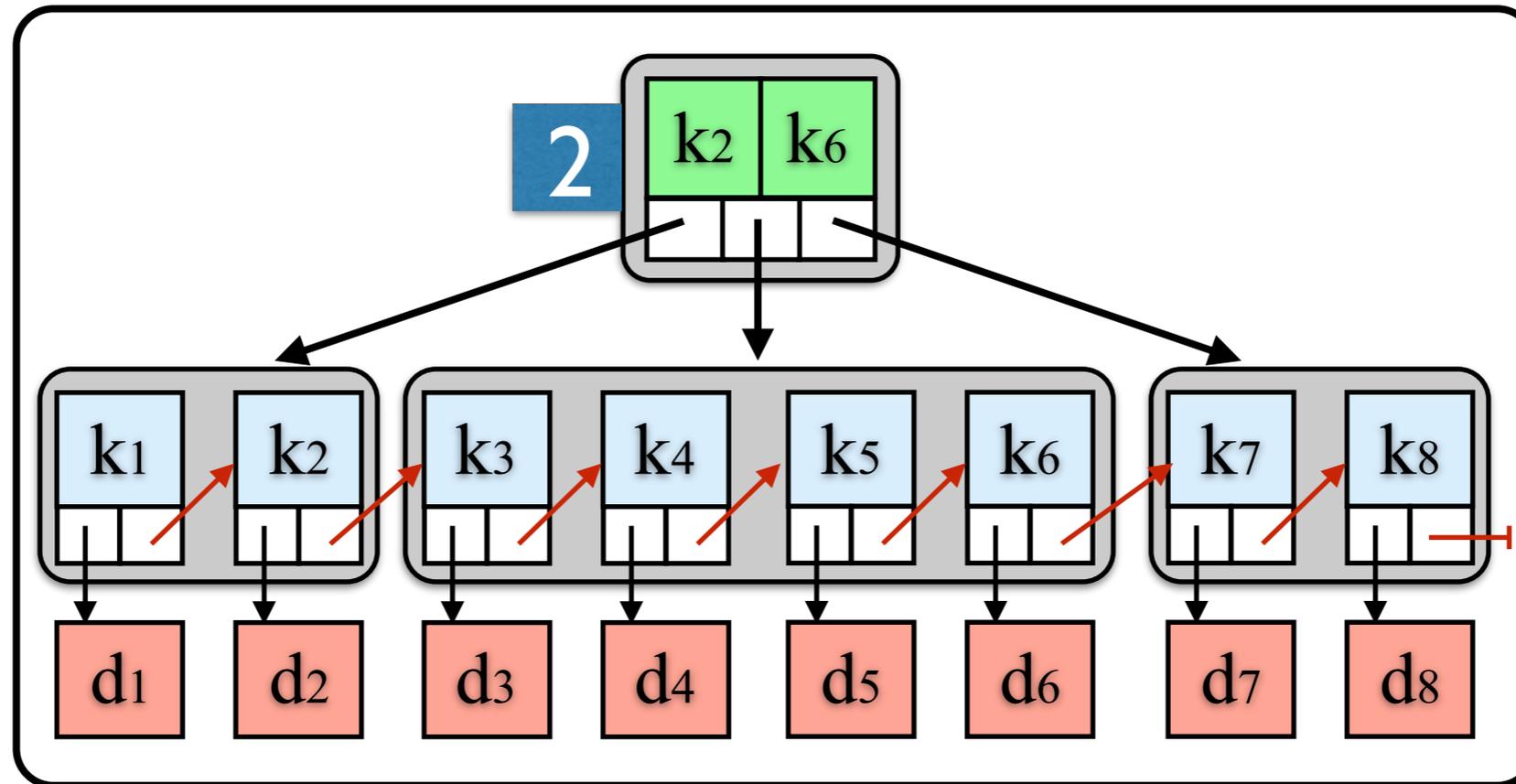
- ❖ Subjective/ Overlapping Shared Resources

- ◆ Proof modularity; better abstraction

- ❖ Local Proofs

- ◆ Proof reuse - proofs done for the largest possible context

B+ Tree $\langle K, V \rangle$: Possible Extension

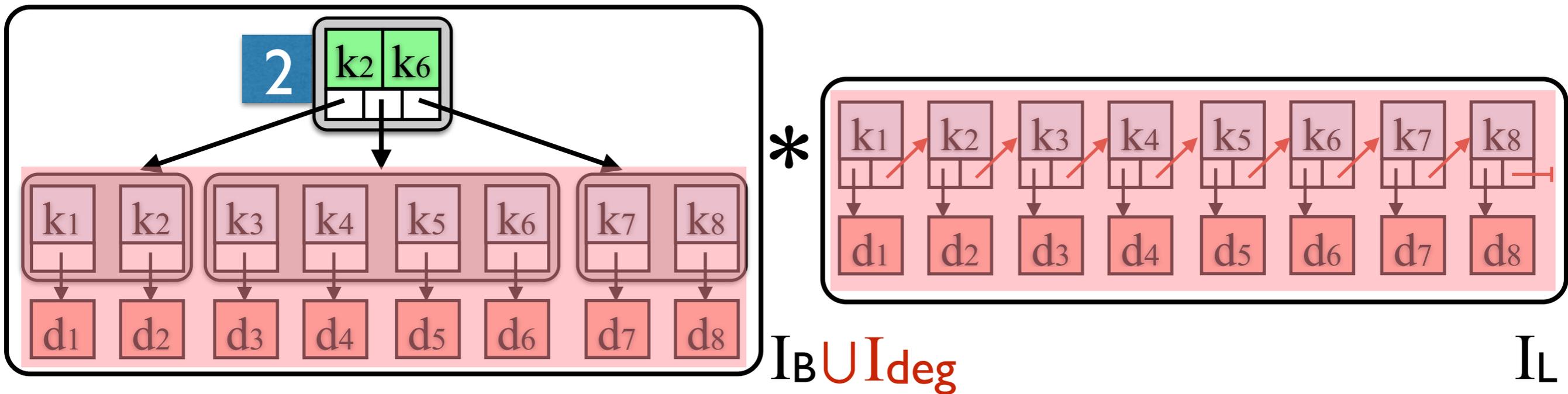
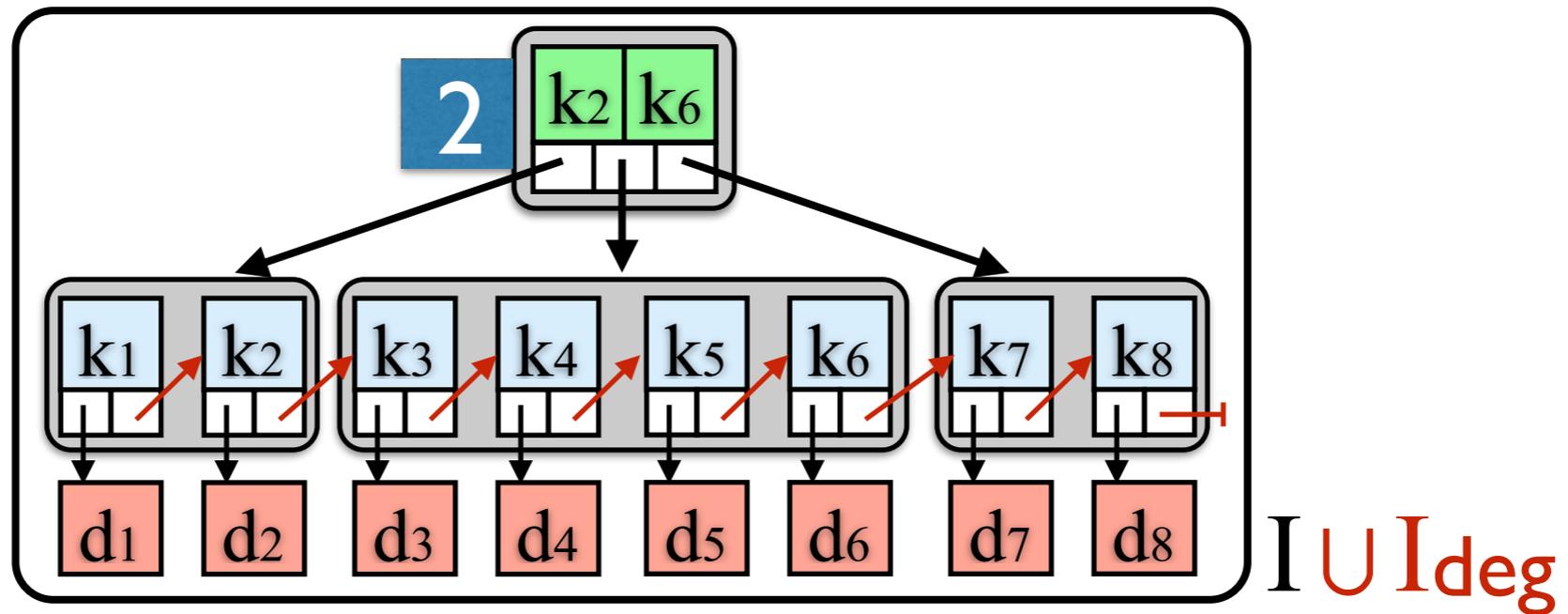


I U Ideg

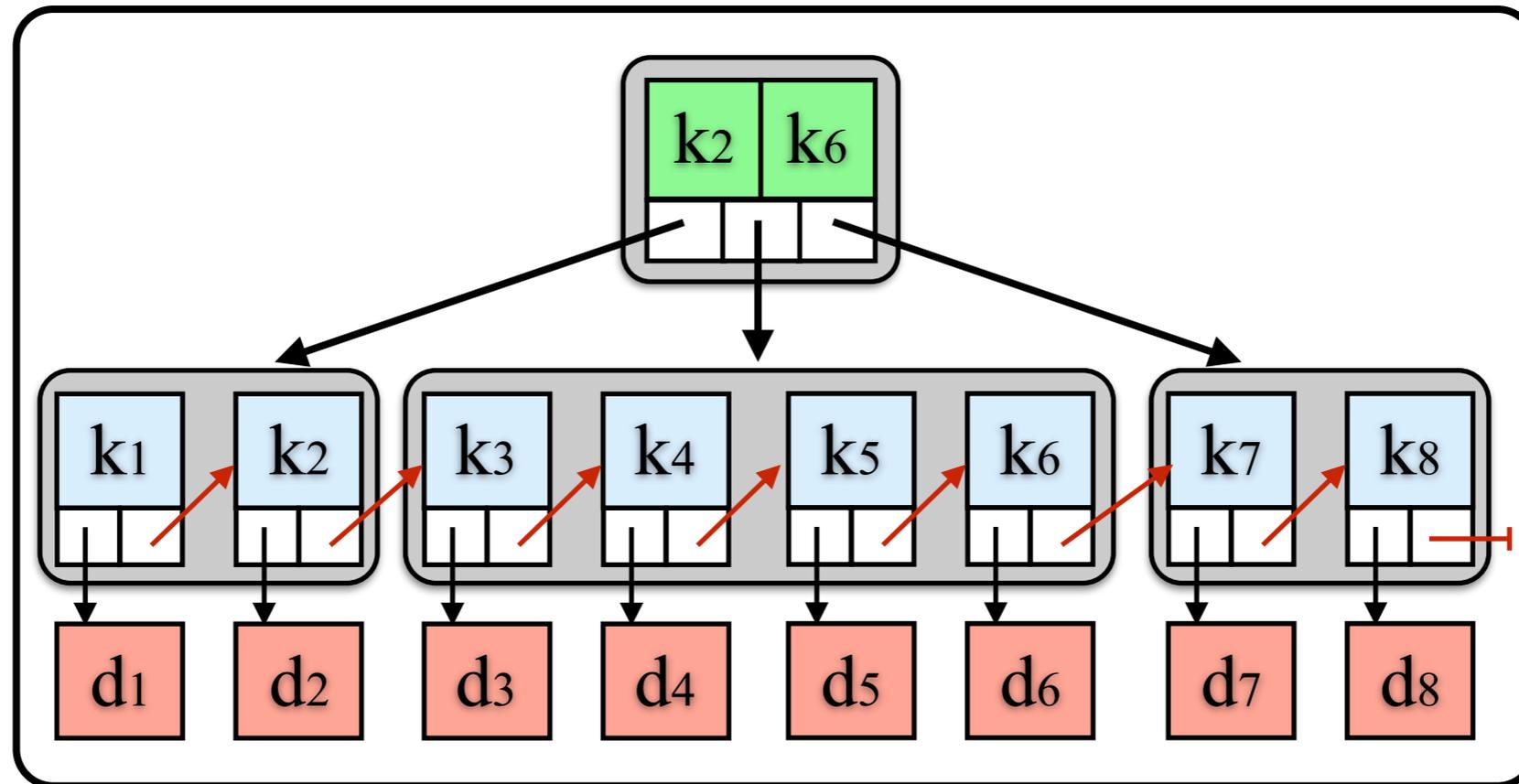
- ❖ Re-degree the tree
 - ✦ periodically change the degree (d) for better search time
- ❖ Re-degreeing ONLY affects the tree structure

Ideg: 2 → 3
- ❖ Re-degreeing does NOT affect the list structure
 - ✦ Should not have to reprove list operations

B+ Tree $\langle K, V \rangle$: Possible Extension (Existing Approaches)



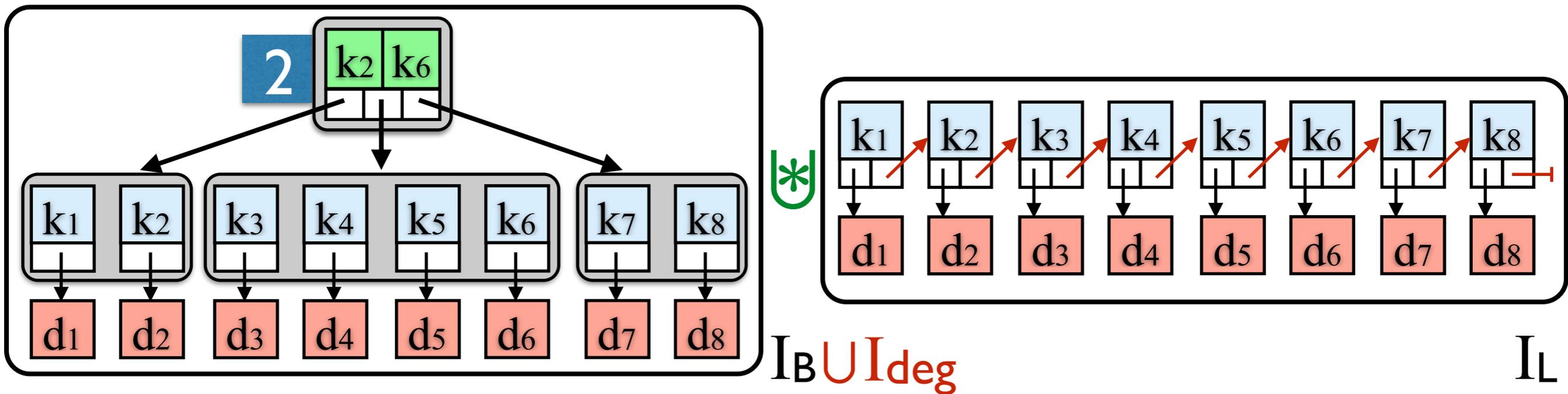
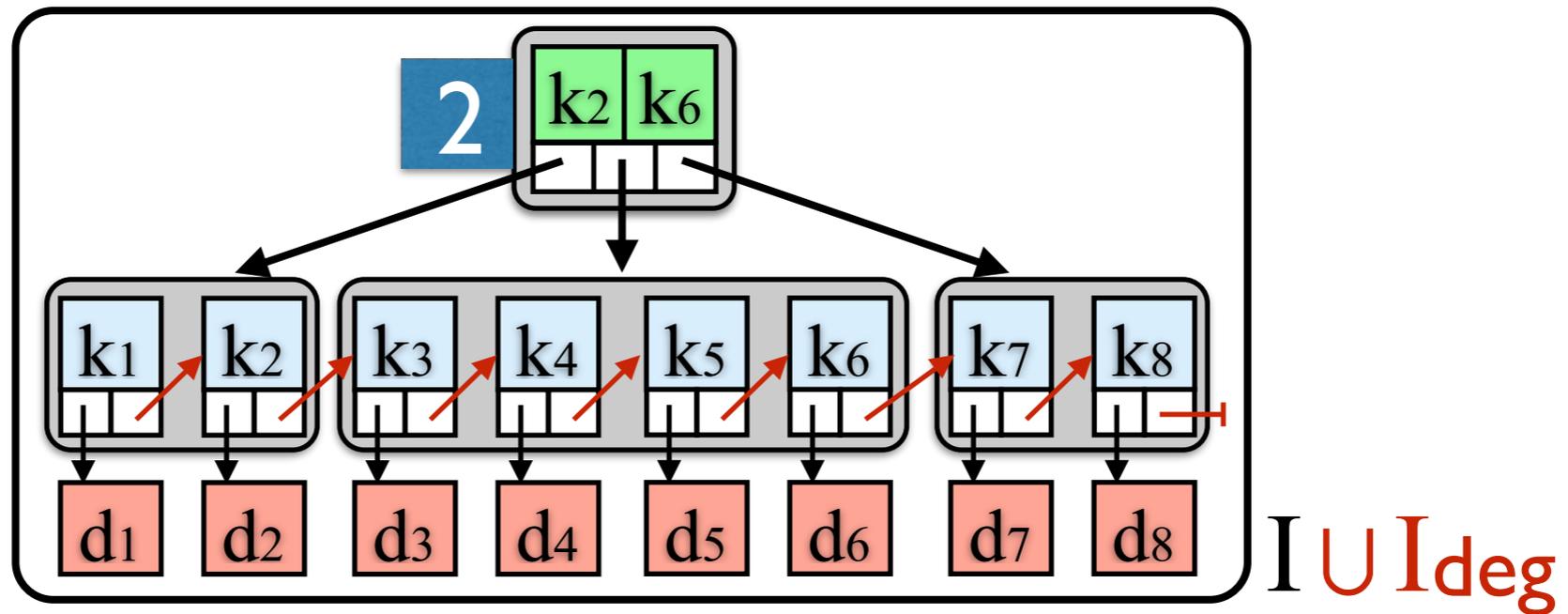
B+ Tree $\langle K, V \rangle$: Possible Extension (Existing Approaches)



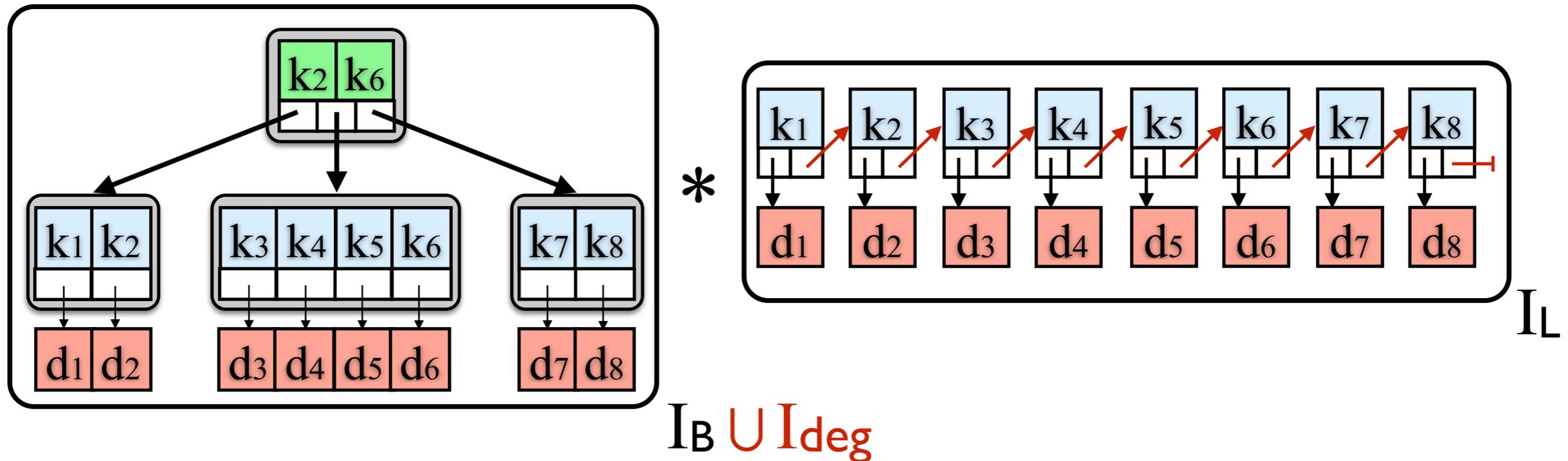
IU Ideg

- ❖ Re-degreeing extension affects ALL proofs
 - ◆ Have to redo all proofs, even those of list-only operations

B+ Tree $\langle K, V \rangle$: Possible Extension (CoLoSL)



B+ Tree $\langle K, V \rangle$: Possible Extension (CoLoSL)



- ❖ Re-degreeing extension only affects the tree
- ❖ Re-degreeing extension does NOT affect the list proofs
 - ◆ Can reuse the proofs as before

Conclusions

- ❖ From OG/Rg to CAP/TaDA
 - ◆ Huge steps towards compositionality/locality
 - ◆ Are we there yet? **No!**
- ❖ CoLoSL
 - ◆ Subjective/overlapping views
 - ◆ Dynamic framing on shared resource/interference
 - ◆ Dynamic extension
 - ◆ Are we there yet? **Still No!**
 - ◆ Abstraction layers; abstract atomicity, ...

Questions?

Thank you for listening